

Algebraic compression of quantum circuits for ~~Hamiltonian~~ simulation

Hermitian



Daan Camps
Lawrence Berkeley National Laboratory
June 16, 2022

Outline and acknowledgements

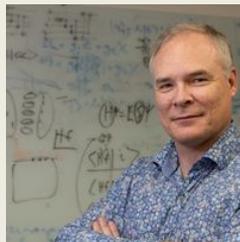
1. A 3 minute **introduction** to quantum computing
2. **Hamiltonian simulation** and Trotter decompositions
3. Algebraic **compression** of Trotterized circuits for spin Hamiltonians
4. **Results** on classical and quantum hardware
5. Conclusion



Efehan Kökcü



**Lindsay
Bassman**



**Bert de
Jong**



**Lex
Kemper**



**Roel Van
Beeumen**

An Algebraic Quantum Circuit Compression Algorithm for Hamiltonian Simulation, D. Camps, E. Kökcü, L. Bassman, W. A. de Jong, A. F. Kemper, R. Van Beeumen, Accepted in SIMAX, arXiv:2108.03283

Algebraic compression of quantum circuits for Hamiltonian evolution, E. Kökcü, D. Camps, L. Bassman, J. K. Freericks, W. A. de Jong, R. Van Beeumen, A. F. Kemper, Phys. Rev. A 105, 032420, arXiv:2108.03282



Introduction to Quantum Computing



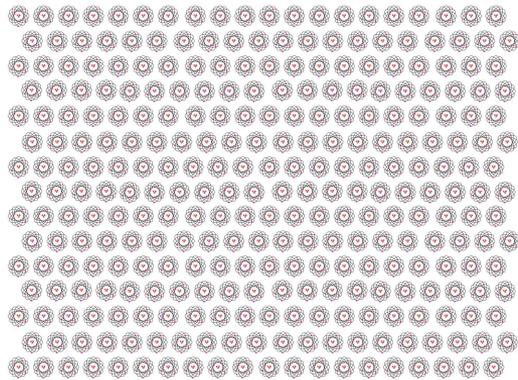
BERKELEY LAB



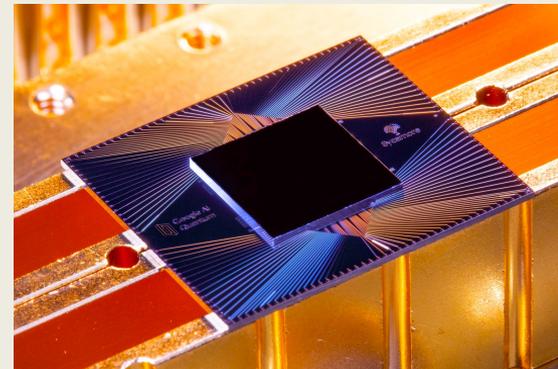
U.S. DEPARTMENT OF
ENERGY

Office of
Science

Dimension of a quantum state grows exponentially with the number of particles



A complete description of a typical quantum state of just 300 qubits requires more bits than the number of atoms in the visible universe (figure from John Preskill).



Google Sycamore chip (2019)

53 qubits

$2^{53} \approx 9 \cdot 10^{15} \approx 36\text{PB}$ (single precision)

$2^{300} =$

2037035976334486086268445688409378161051468393665936250636140449354381299763336706183397376

Quantum computing from 10000 ft

Two things are required for quantum computation:

- An **encoding** of the data in the quantum state $|\Psi\rangle$
- A way to **control** the evolution towards an encoding of the solution

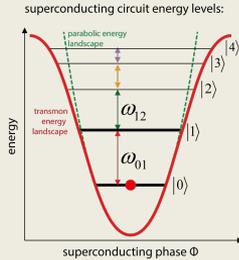
Quantum computers
are coherently
controllable quantum
systems



Advanced Quantum
Testbed @ Berkeley Lab

Qubits represent quantum data

Physics: two-level quantum system



Math: 2-dimensional complex vectors with unit norm

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$$

$$|\alpha|^2 + |\beta|^2 = 1$$

Quantum Gates: change state of a qubit

$$|\psi\rangle \rightarrow \boxed{U} |\psi\rangle \rightarrow |\psi'\rangle = U|\psi\rangle$$

U is **unitary**

Unitary matrices preserve the norm of the vector (quantum operations are Hamiltonian time evolution)

Multi-qubit states and quantum circuits

$$|00\rangle = |0\rangle \otimes |0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \otimes \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$|01\rangle = |0\rangle \otimes |1\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \otimes \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}$$

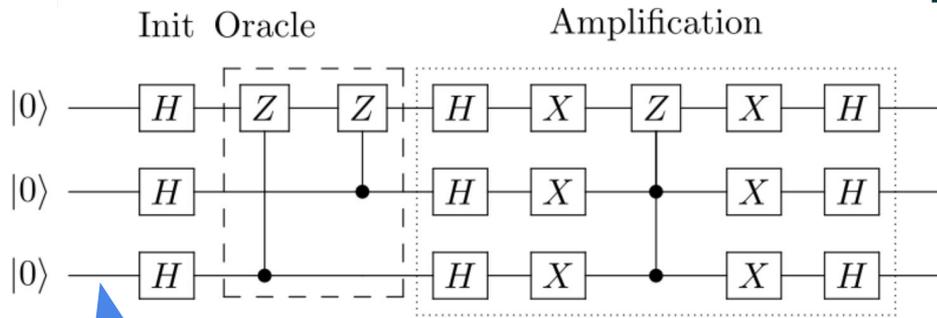
$$|10\rangle = |1\rangle \otimes |0\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \otimes \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}$$

$$|11\rangle = |1\rangle \otimes |1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \otimes \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

1 qubit → 2 basis states
 2 qubits → 4 basis states

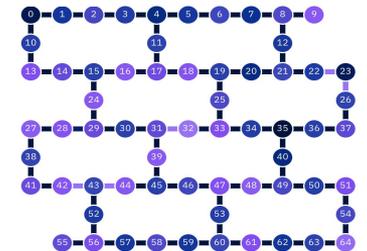
 n qubits → 2^n basis states

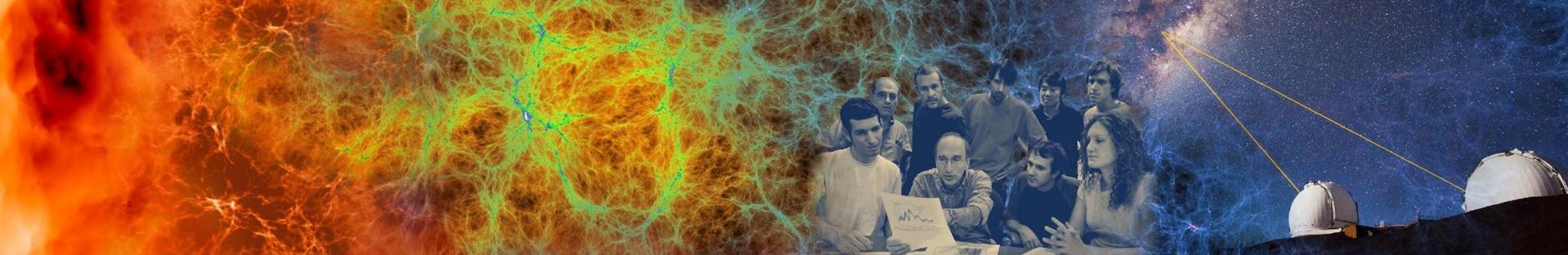
$$|\psi\rangle = \alpha_0 |00 \dots 0\rangle + \alpha_1 |00 \dots 1\rangle + \dots + \alpha_{2^n-1} |11 \dots 1\rangle = \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \vdots \\ \alpha_{2^n-1} \end{bmatrix} \in \mathbb{C}^{2^n}$$



each line is a qubit

direction of computation





Hamiltonian simulation and Trotterization

Hamiltonian simulation

Simulate time evolution under Schrödinger equation for a time-dependent Hamiltonian

$$\frac{\partial}{\partial t} \psi(t) = -iH(t)\psi(t) \quad H(t) \in \mathbb{C}^{2^N \times 2^N}$$

Hermitian matrix

Solved by applying the time-evolution operator:

$$U(t_1, t_0) = \mathcal{T} \exp \left(-i \int_{t_0}^{t_1} H(t) dt \right)$$

to the initial state:
$$\psi(t_1) = U(t_1, t_0)\psi(t_0)$$

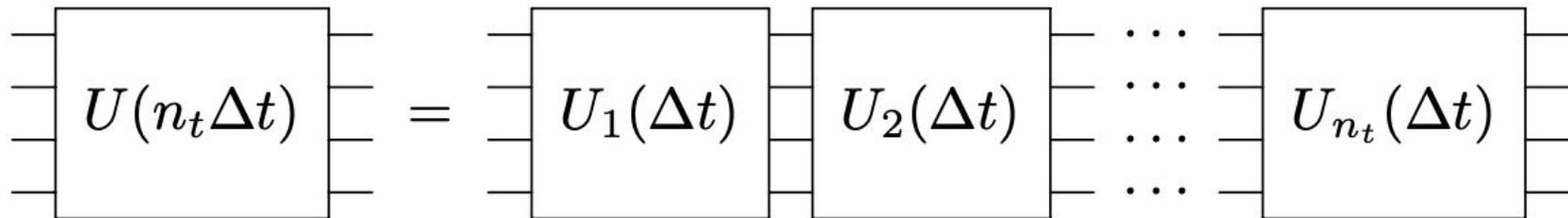
Time-independent case:
$$U(t_1, t_0) = \exp(-i(t_1 - t_0)H)$$

Trotter splitting and time discretization

Trotter decomposition (or operator splitting):

$$H = A + B \quad U(\Delta t) = \exp(-iA\Delta t) \exp(-iB\Delta t)$$

$$\|U(\Delta t) - \exp(-iH\Delta t)\| \leq \frac{\Delta t^2}{2} \|[A, B]\|$$



1D Spin- $1/2$ Hamiltonians

Pauli spin- $1/2$ matrices: $\sigma^x = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$, $\sigma^y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}$, $\sigma^z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$

Basis for: $\mathfrak{su}(2)$ Generators for: $SU(2) = \left\{ \begin{bmatrix} \alpha & -\bar{\beta} \\ \beta & \bar{\alpha} \end{bmatrix} : \alpha, \beta \in \mathbb{C}, |\alpha|^2 + |\beta|^2 = 1 \right\}$

$$\sigma_i^\alpha := \underbrace{I \otimes \dots \otimes I}_{i-1} \otimes \sigma^\alpha \otimes \underbrace{I \otimes \dots \otimes I}_{N-i}$$

Transverse field XY model:

$$H(t) = \underbrace{\sum_{i=1}^{N-1} J_i^x(t) \sigma_i^x \sigma_{i+1}^x + J_i^y(t) \sigma_i^y \sigma_{i+1}^y}_{\text{Coupling}} + \underbrace{\sum_{i=1}^N h_i^z(t) \sigma_i^z}_{\text{External Field}}$$

Circuit diagrams

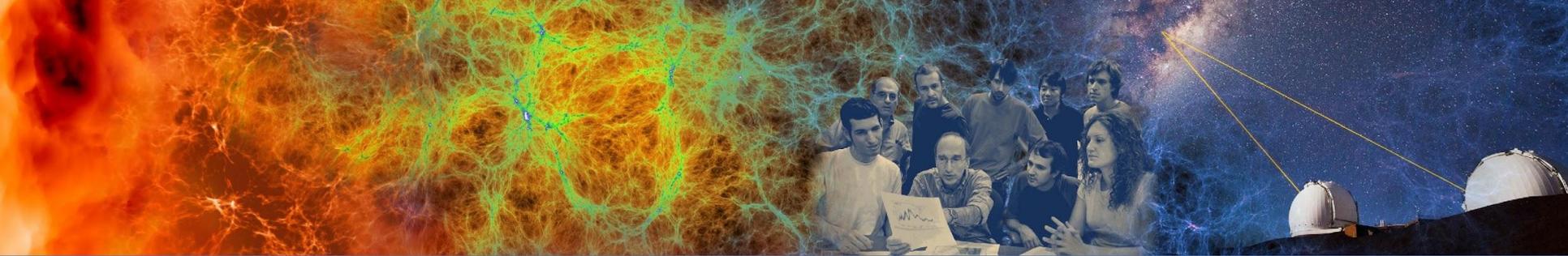
Single-qubit rotation over Pauli- α axis ($\alpha \in \{x, y, z\}$):

$$R^\alpha(\theta) := \exp(-i \sigma^\alpha \theta/2) = \text{---} \boxed{\alpha} \text{---}$$

Two-qubit rotation over Pauli- α axis ($\alpha \in \{x, y, z\}$):

$$R^{\alpha\alpha}(\theta) := \exp(-i \sigma^\alpha \otimes \sigma^\alpha \theta/2) = \text{---} \boxed{\alpha} \text{---}$$

Easy operations to execute on QC: Native gate for ion traps, 2 CNOTs for superconducting



Algebraic compression of Hamiltonian simulation circuits

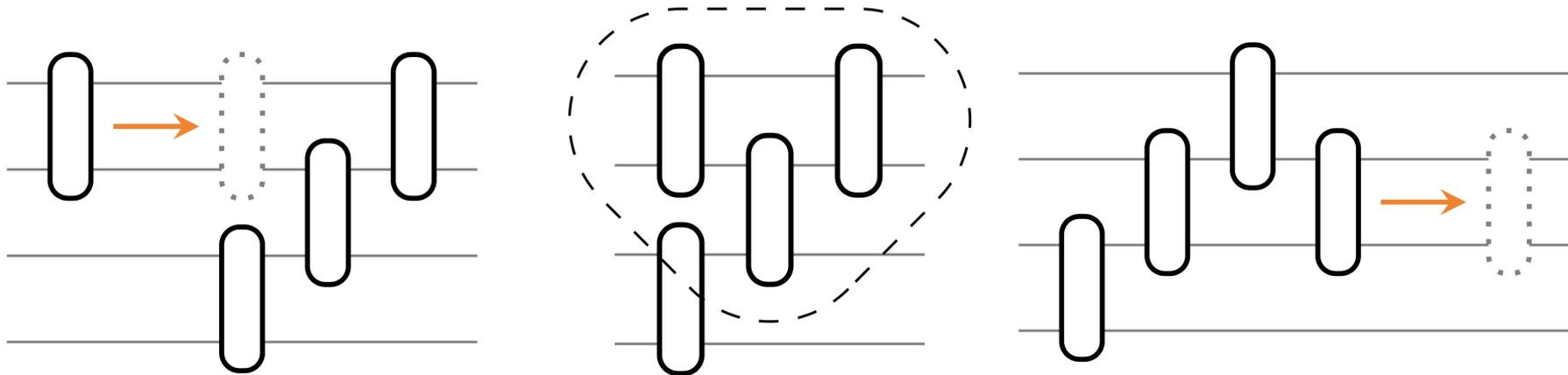
Definition

A **block** is a parametrized and indexed family of operators $\mathbf{B}_i(\boldsymbol{\theta})$ that satisfy 3 properties:

- **Fusion:**

$$\mathbf{B}_i(\boldsymbol{\theta}_1)\mathbf{B}_i(\boldsymbol{\theta}_2) = \mathbf{B}_i(\hat{\boldsymbol{\theta}}) \quad i \quad i = i$$

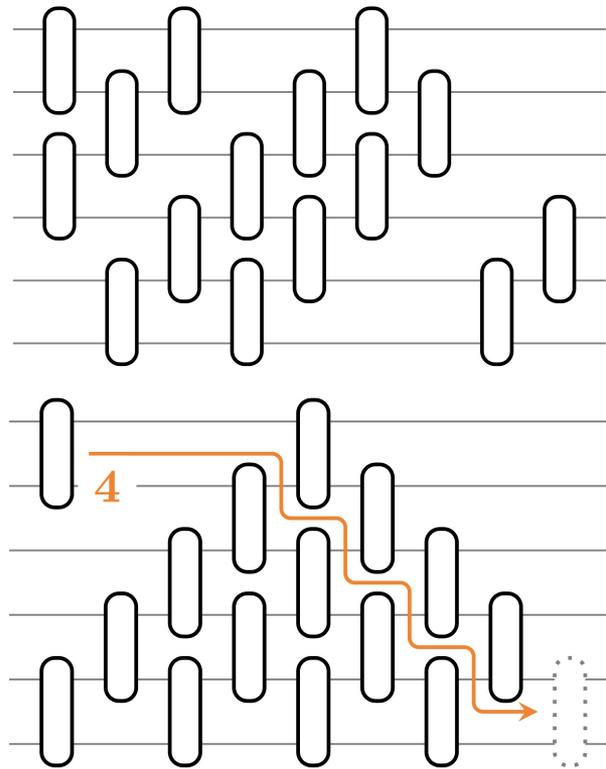
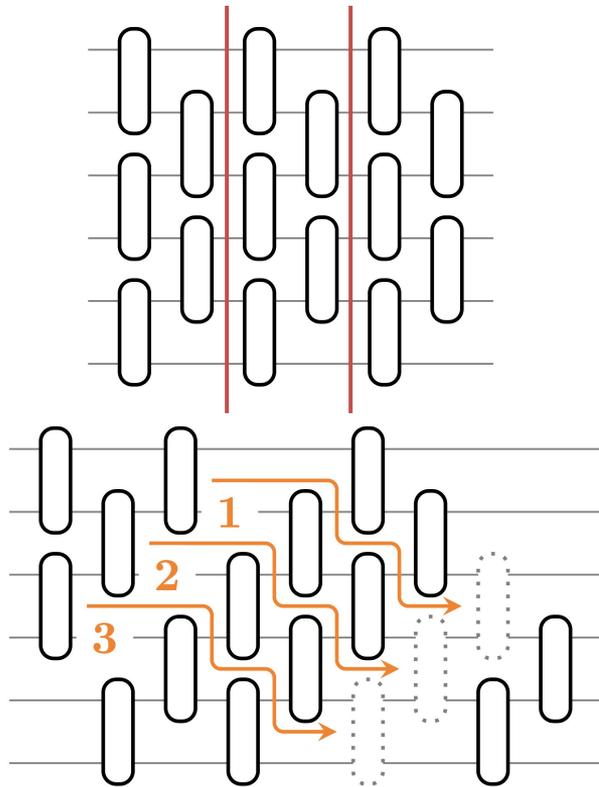
Central mechanism in our compression algorithm



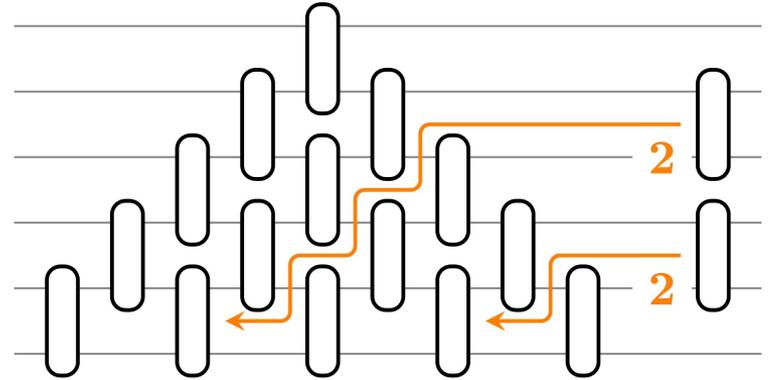
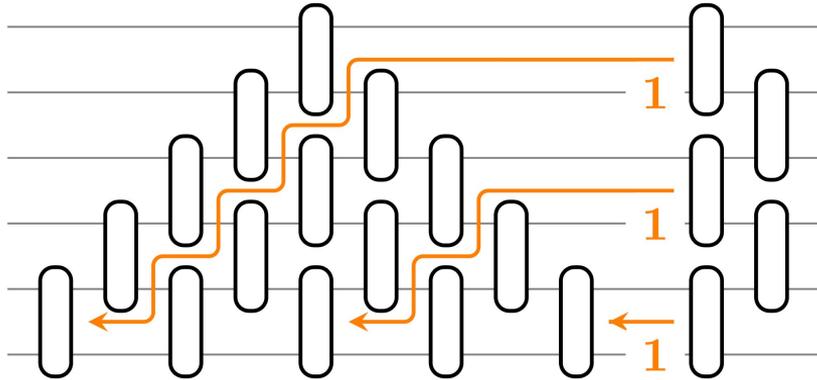
Equivalent mechanism as in *core-chasing* eigenvalue algorithms, but on operators of exponential dimension.

Core-Chasing Algorithms for the Eigenvalue Problem, Aurentz, Mach, Robol, Vandebril, Watkins

Transforming squares to triangles



Merging time-steps into triangles



Even/odd blocks act on independent parts of the triangle. In our implementation, these are merged in parallel

Euler decomposition and turnover of SU(2)

Lemma: Euler decomposition

Let $\alpha, \beta \in \{x, y, z\}, \alpha \neq \beta$. Every $U \in \text{SU}(2)$ can be represented as:

$$U = R^\alpha(\theta_1) R^\beta(\theta_2) R^\alpha(\theta_3), \quad \text{---} \textcircled{U} \text{---} = \text{---} \textcircled{\alpha}_{\theta_3} \textcircled{\beta}_{\theta_2} \textcircled{\alpha}_{\theta_1} \text{---}$$

Lemma: SU(2) turnover

Let $\alpha, \beta \in \{x, y, z\}, \alpha \neq \beta$. For every $\theta_1, \theta_2, \theta_3$, there exist $\theta_a, \theta_b, \theta_c$ such that

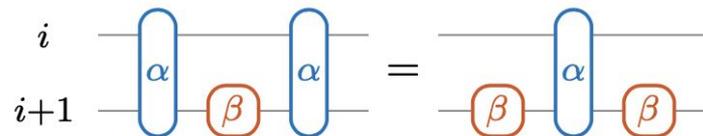
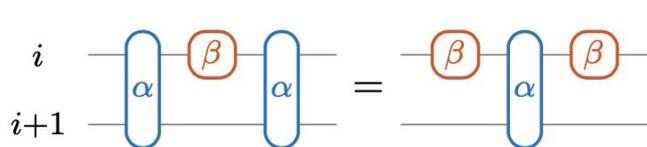
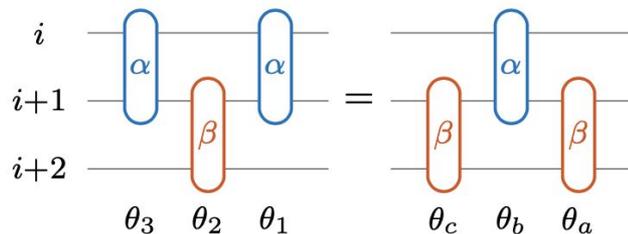
$$R^\alpha(\theta_1) R^\beta(\theta_2) R^\alpha(\theta_3) = R^\beta(\theta_a) R^\alpha(\theta_b) R^\beta(\theta_c), \quad \text{---} \textcircled{\alpha}_{\theta_3} \textcircled{\beta}_{\theta_2} \textcircled{\alpha}_{\theta_1} \text{---} = \text{---} \textcircled{\beta}_{\theta_c} \textcircled{\alpha}_{\theta_b} \textcircled{\beta}_{\theta_a} \text{---}$$

→ We can compute the SU(2) turnover backward stable (Givens rotations)

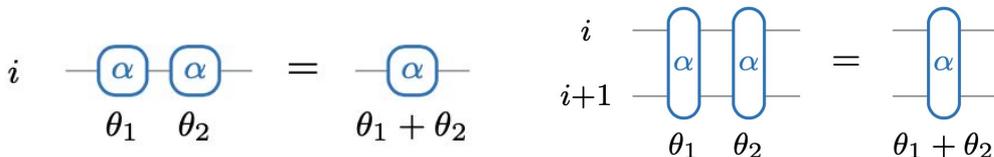
SU(2) groups in disguise

Lemma:

Let $\alpha, \beta \in \{x, y, z\}, \alpha \neq \beta$. The following operations are also dual Euler decompositions of SU(2):



Fusion operations are trivial:

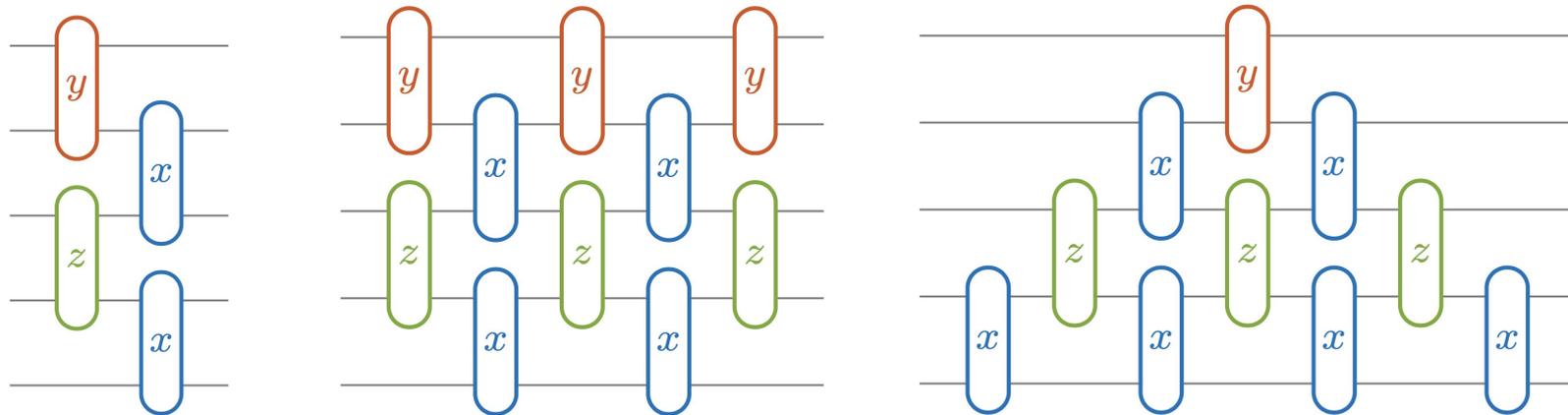


Kitaev Chain

A Kitaev chain is a Hamiltonian of the form:

$$H(t) = \sum_{i=1}^{N-1} J_i^{\alpha_i}(t) \sigma_i^{\alpha_i} \sigma_{i+1}^{\alpha_i} \quad \alpha_i \neq \alpha_{i+1}$$

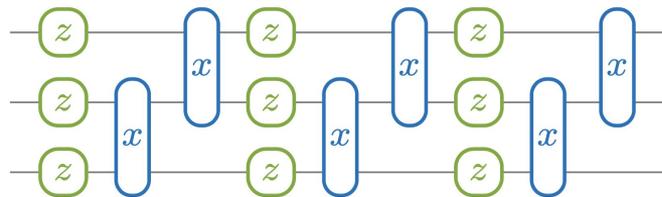
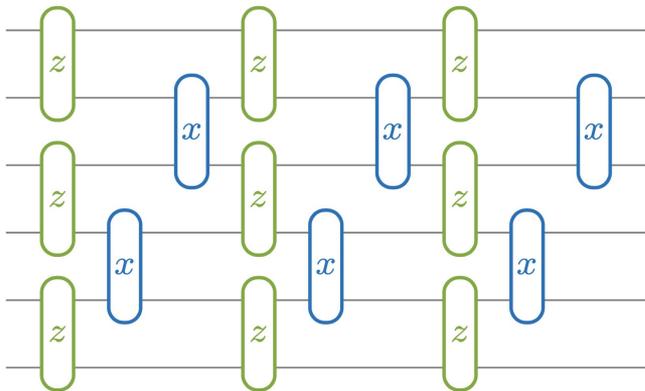
For example: $H(t) = J_1^y(t) \sigma_1^y \sigma_2^y + J_2^x(t) \sigma_2^x \sigma_3^x + J_3^z(t) \sigma_3^z \sigma_4^z + J_4^x(t) \sigma_4^x \sigma_5^x$



TFIM Hamiltonian

The Transverse-Field Ising Model has the form:

$$H(t) = \sum_{i=1}^{N-1} J_i^\alpha(t) \sigma_i^\alpha \sigma_{i+1}^\alpha + \sum_{i=1}^N h_i^\beta(t) \sigma_i^\beta \quad \alpha \neq \beta$$



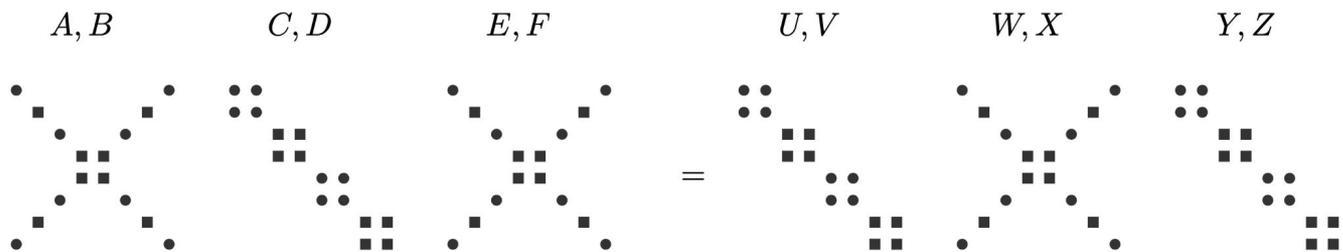
N-qubit TFIM is isomorphic to 2N-qubit Kitaev chain

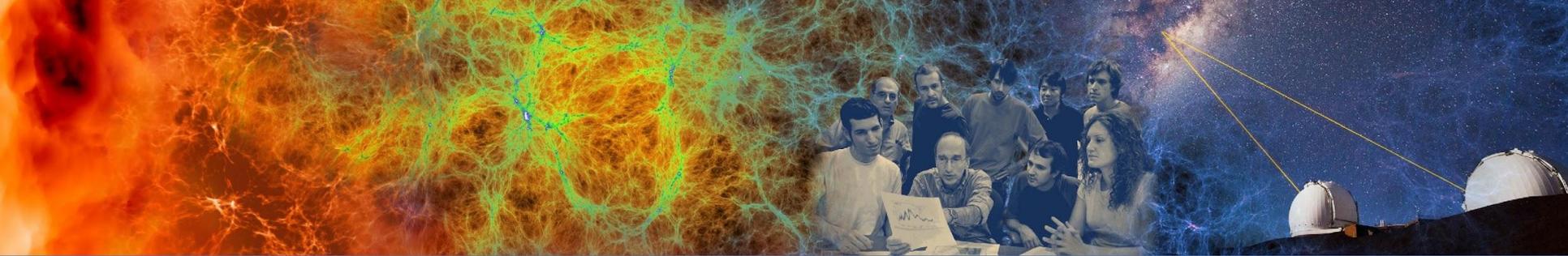
TFXY Hamiltonian

Transverse field XY model:
$$H(t) = \underbrace{\sum_{i=1}^{N-1} J_i^x(t) \sigma_i^x \sigma_{i+1}^x + J_i^y(t) \sigma_i^y \sigma_{i+1}^y}_{\text{Coupling}} + \underbrace{\sum_{i=1}^N h_i^z(t) \sigma_i^z}_{\text{External Field}}$$

TFXY block:
$$= \begin{bmatrix} \alpha & & & -\bar{\delta} \\ & \beta & -\bar{\gamma} & \\ & \gamma & \bar{\beta} & \\ \delta & & & \bar{\alpha} \end{bmatrix}$$

Turnover through simultaneous diagonalization





Results



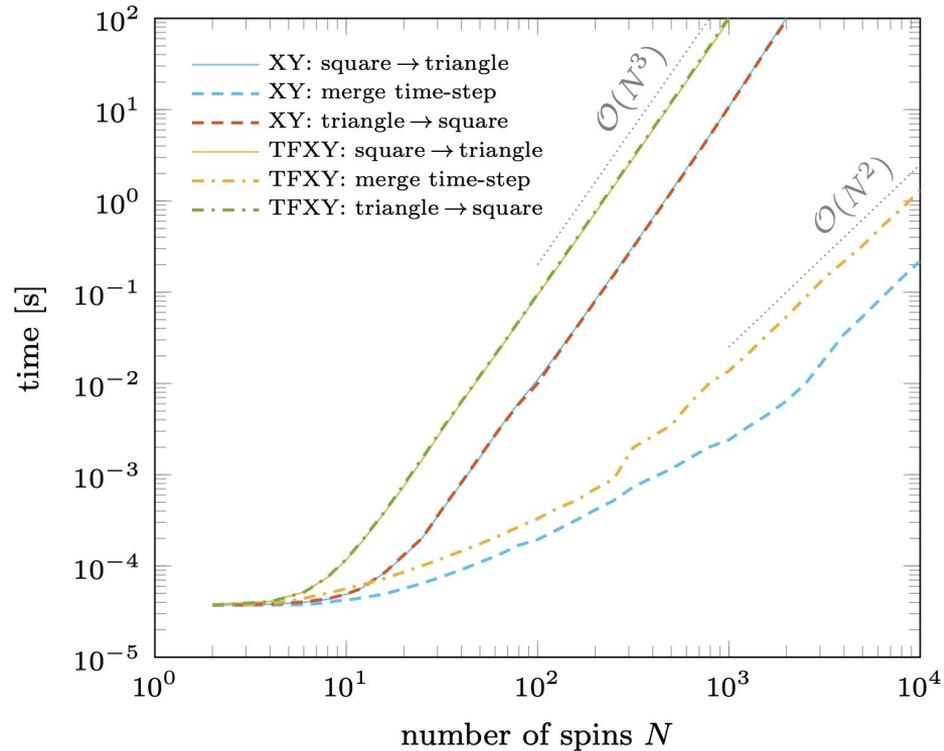
BERKELEY LAB



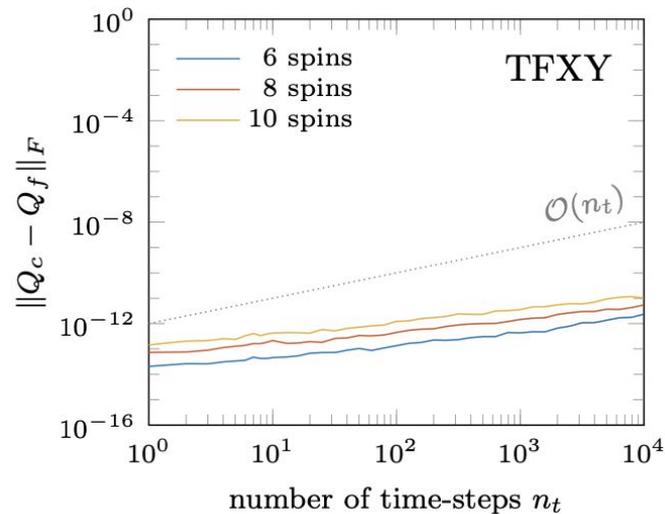
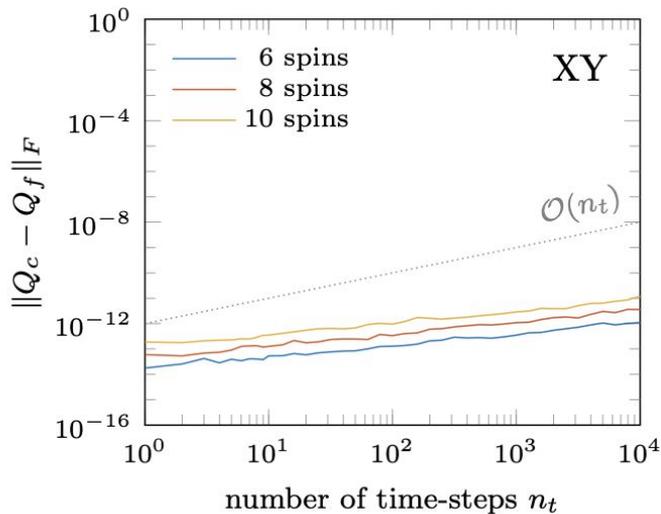
U.S. DEPARTMENT OF
ENERGY

Office of
Science

Numerical results: timings



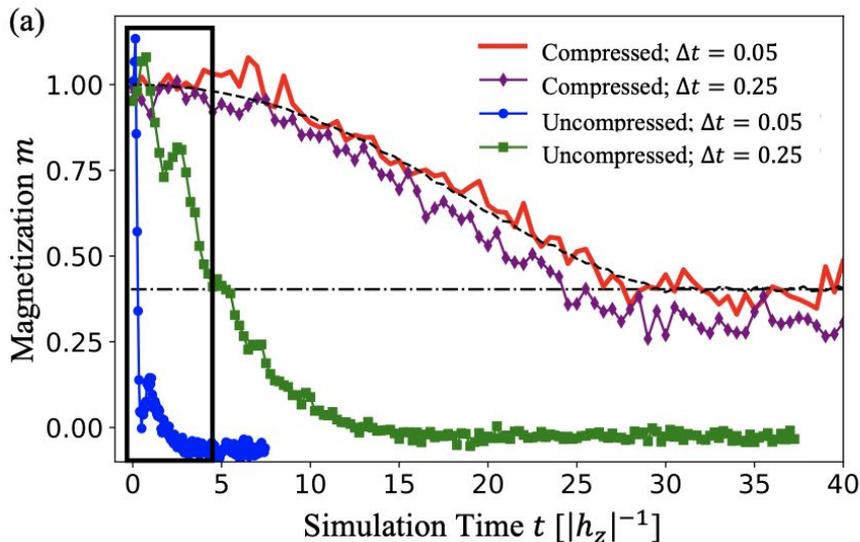
Numerical results: backward error



Quantum Computer: Adiabatic State Preparation

- Time evolve TFIM in ground state from trivial state w/o coupling to more complicated ground state with coupling terms
- 5 qubit model on IBMQ Brooklyn
- Measure the average magnetization

$$H(t) = J(t) \sum_i \sigma_i^x \sigma_{i+1}^x + h \sum_i \sigma_i^z$$



Conclusion

- Efficient and stable **classical** numerical algorithm for compression of quantum circuits for simulation of **integrable TFX chains**
- **Enables simulation** of small systems on current generation noisy quantum hardware
 - Prepare **non-trivial states**
 - Simulate interesting physics phenomena
- Extensions to 2D non-interacting, controlled evolutions, ...

Fast Free Fermion Compiler (F3C):
<https://github.com/QuantumComputingLab>

arXiv:2108.03282, arXiv:2108.03283

