

Pole swapping methods for the eigenvalue problem

Rational QR algorithms

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Overview

Introduction

- Generalized eigenvalue problems
- Bulge chasing
- Pole swapping
 - Rational Krylov
 - Rational QZ
 - Rational accelerated subspace iteration
- Multishift, multipole rational QZ
- Conclusion

Introduction

- Let $A, B \in \mathbb{F}^{n \times n}$ define a matrix pair (A, B) or matrix pencil $A \lambda B$.
- Regular: $det(A \lambda B) \neq 0$.
- Generalized eigenvalue problem:

$$A\mathbf{x} = \lambda B\mathbf{x}, \quad \lambda \in \overline{\mathbb{C}} := \mathbb{C} \cup \{\infty\}.$$

• For a regular matrix pencil $A - \lambda B$ there exists unitary matrices Q and Z such that

 $S - \lambda T = Q^* (A - \lambda B) Z$

with $S - \lambda T$ upper triangular and $\Lambda(A, B) = \{s_{11}/t_{11}, s_{22}/t_{22}, \ldots\}$.

• The QZ method (Moler-Stewart, 1973) is the default algorithm to compute the generalized Schur decomposition.

- The QZ method generalizes the QR method of (Francis, 1961-62) for the standard eigenvalue problem
- It is a bulge chasing algorithm which consists out of two phases:
 - 1. Initial (direct) reduction to equivalent Hessenberg, upper triangular form

 $H - \lambda R = Q^* (A - \lambda B) Z$

2. Iterative bulge chasing phase to compute (real) generalized Schur decomposition

 $S - \lambda T = Q^* (A - \lambda B) Z$

Bulge chasing =

• Motivated by implicit Q theorems

 \Rightarrow iterates are uniquely determined by $\boldsymbol{q}_1 = p(AB^{-1})\boldsymbol{e}_1$

• Nested subspace iteration with a change of basis accelerated by polynomials (shifts) (Elsner-Watkins, 1991; Watkins, 1993)

 \rightarrow These results are based on a connection with Krylov subspaces.

Pole swapping

| Polynomial | Rational | | |
|---|--|--|--|
| | rational Krylov subspace | | |
| $\frac{Krylov subspace}{\mathcal{K}_{m+1}(A, \mathbf{v}) := \mathcal{R}(\mathbf{v}, A\mathbf{v}, \dots, A^m\mathbf{v})$ | $\mathcal{K}^{rat}_{m+1}(A,oldsymbol{ u},\Xi):=q(A)^{-1}\mathcal{K}_{m+1}(A,oldsymbol{ u})$ | | |
| | $\bullet \equiv = (\xi_1, \ldots, \xi_m) \subset \bar{\mathbb{C}} \setminus \Lambda$ | | |
| | • $q(z) = \prod_{\xi_i eq \infty} (z - \xi_i)$ | | |
| $\frac{\text{Arnoldi Decomposition}}{AV_m = V_{m+1}\underline{H}_m}$ (Arnoldi, 1951) | rational Arnoldi (Ruhe, 1998) | | |
| | $AV_{m+1}\underline{K}_m = V_{m+1}\underline{L}_m$ | | |
| | • $\ell_{i+1,i}/k_{i+1,i} = \xi_i$ | | |
| Uniqueness | Uniqueness (Berljafa-Güttel, 2015) | | |
| $m{v}$ fixed $\Leftrightarrow (V_{m+1}, \underline{H}_m)$ unique | $m{ u}$ and Ξ fixed \Leftrightarrow $(V_{m+1}, \underline{K}_m, \underline{L}_m)$ unique | | |

Rational Krylov

Berljafa-Güttel (2015)

Two methods to change the poles in a rational Arnoldi decomposition

$$AV_{m+1}\underline{K}_m = V_{m+1}\underline{L}_m$$

- Implicitly by changing v_1
- Explicitly by pole swapping
- \Rightarrow Equivalent if $\mathcal{R}(V_{m+1})$ is fixed.

C.-Meerbergen-Vandebril (2019a)

The pole swapping technique can be used as a direct method for the eigenvalue problem: rational QZ algorithm.

Rational QZ

Hessenberg pencils



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Rational QZ

Hessenberg pencils

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Rational QZ

Hessenberg pencils



Introducing a shift



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Introducing a shift



Introducing a shift



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Swapping poles



Swapping poles

Classical problem in NLA: Reordering generalized Schur form

(Van Dooren, 1981), (Kågström, 1993), (C.-Mach-Vandebril-Watkins, 2019)

$$\Rightarrow Q^* = \begin{bmatrix} \times & \times \\ \times & \times \end{bmatrix} \quad , Z = \begin{bmatrix} \times & \times \\ \times & \times \end{bmatrix}$$

Swapping poles

Table 1: Distribution of errors $\hat{a}_{21}/||A||$ and $\hat{b}_{21}/||B||$ for our method, Van Dooren's method, and the Sylvester method.

| $\hat{x}_{21}/\ X\ $ | | $\left[0,10^{-16}\right]$ | $\left(10^{-16},10^{-15}\right]$ | $\left(10^{-15},10^{-10}\right]$ | $\left(10^{-10},10^{-5}\right]$ | $\left(10^{-5},10^0\right]$ |
|----------------------|---|---------------------------|----------------------------------|----------------------------------|---------------------------------|-----------------------------|
| Our method | A | 99.71% | 0.29% | 0% | 0% | 0% |
| | B | 99.85% | 0.15% | 0% | 0% | 0% |
| Van Dooren | A | 98.19% | 0.55% | 0.93% | 0.27% | 0.06% |
| | B | 98.19% | 0.55% | 0.93% | 0.27% | 0.06% |
| Sylvester | A | 93.34% | 5.88% | 0.57% | 0.17% | 0.04% |
| | B | 93.34% | 5.88% | 0.57% | 0.17% | 0.04% |

Swapping poles



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Swapping poles



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Swapping poles





Swapping poles



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Swapping poles







Swapping poles









Swapping poles



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Introducing a pole



Introducing a pole

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Classical QZ as a special case

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Classical QZ as a special case

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Classical QZ as a special case

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Definition: Properness.

The Hessenberg pair (A, B) is called *proper* if:







Theorem. (C.-Meerbergen-Vandebril, 2019a)

If (A, B) is a proper Hessenberg pair with poles $(\xi_1, \ldots, \xi_{n-1})$ distinct from the eigenvalues. Then for $i = 1, \ldots, n$:

$$\mathcal{K}_i^{\mathsf{rat}}(AB^{-1}, \boldsymbol{e}_1, (\xi_1, \ldots, \xi_{i-1})) = \mathcal{E}_i := \mathcal{R}(\boldsymbol{e}_1, \ldots, \boldsymbol{e}_i),$$

while for i = 1, ..., n - 1:

 $\mathcal{K}_i^{\mathsf{rat}}(B^{-1}A, \boldsymbol{e}_1, (\xi_2, \ldots, \xi_i)) = \mathcal{E}_i.$

Implicit Q Theorem. (C.-Meerbergen-Vandebril, 2019a)

Given a regular matrix pair (A, B). The matrices Q and Z that transform it to proper Hessenberg form,

$$(\hat{A},\hat{B})=Q^*(A,B) Z,$$

are determined essentially unique if Qe_1 and the (order of the) poles are fixed.

Theoretical results

Rational accelerated subspace iteration. (C.-Meerbergen-Vandebril, 2019a) A rational QZ step with shift $\varrho \notin \{\Lambda, \Xi\}$ on a proper Hessenberg pencil with poles $(\xi_1, \ldots, \xi_{n-1})$ and new pole ξ_n , all distinct from Λ , performs nested subspace iteration for $i = 1, \ldots, n-1$ accelerated by

$$Q\mathcal{E}_i = \mathcal{R}(\boldsymbol{q}_1, \dots, \boldsymbol{q}_i) = (A - \varrho B)(A - \xi_i B)^{-1} \mathcal{E}_i$$
$$Z\mathcal{E}_i = \mathcal{R}(\boldsymbol{z}_1, \dots, \boldsymbol{z}_i) = (A - \xi_{i+1}B)^{-1}(A - \varrho B)\mathcal{E}_i.$$

followed by a change of basis.

 \rightarrow Subspace iteration with rational filter \rightarrow More modular (single swap) convergence theory: (C.-Mach-Vandebril-Watkins, 2019).

Exactness result (C., 2019)

Let (A, B) be a proper Hessenberg pencil with poles Ξ . Furthermore, let ρ be an eigenvalue of (A, B) which is distinct from Ξ . A rational QZ step, $Q^*(A, B)Z$, with shift ρ leads to a deflation in the last rows of $Q^*(A, B)Z$.

$\mathsf{Pole}\ \mathsf{swapping} =$

• Motivated by implicit Q theorems

 \Rightarrow iterates are uniquely determined by $m{q}_1=q(AB^{-1})m{e}_1$ and poles in pencil

- Nested subspace iteration with a change of basis accelerated by rational functions (shifts and poles)
- \rightarrow These results are based on a connection with rational Krylov subspaces

Numerical example: Reduction to Hessenberg form

Data: MHD matrix pair from MatrixMarket, n = 1280



Numerical example: Reduction to Hessenberg form

Data: MHD matrix pair from MatrixMarket, n = 1280



Numerical example: Reduction to Hessenberg form

Data: MHD matrix pair from MatrixMarket, n = 1280



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Motivation: make the rational QZ method competitive with state-of-the-art.

- Extension of the rational QZ method from Hessenberg to block Hessenberg pencils
- Shifts and poles of larger multiplicity
- Real-valued generalized eigenproblems in real arithmetic
- Swapping 2×2 blocks: Iterative refinement via Newton steps (C.-Mastronardi-Vandebril-Van Dooren, 2019).



Aggressive early deflation (Braman-Byers-Mathias, 2002)



Numerical experiments with libRQZ v0.1



Conclusion

1. We have presented a novel interpretation of QR-type methods:

bulge chasing \leftrightarrow pole swapping.

- 2. This results in a more general class of algorithms.
- 3. Convergence is determined by rational functions instead of polynomials.
- 4. Faster and more flexible eigensolvers.

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