

RQZ: A rational QZ method for the generalized eigenvalue problem

Daan Camps, Karl Meerbergen, and Raf Vandebril

May 6, 2018

KU Leuven - University of Leuven - Department of Computer Science - NUMA Section

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem:

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem:

- ◊ Small to medium-sized
- ◊ Regular
- ◊ All eigenvalues required
- ◊ No symmetry assumed

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem:

- ◊ Small to medium-sized
- ◊ Regular
- ◊ All eigenvalues required
- ◊ No symmetry assumed

Tools:

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem:

- ◊ Small to medium-sized
- ◊ Regular
- ◊ All eigenvalues required
- ◊ No symmetry assumed

Tools:

- ◊ Hessenberg, Hessenberg form
- ◊ Shift & pole introduction and swapping
- ◊ Rational Krylov
- ◊ Subspace iteration

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem:

- ◊ Small to medium-sized
- ◊ Regular
- ◊ All eigenvalues required
- ◊ No symmetry assumed

Tools:

- ◊ Hessenberg, Hessenberg form
- ◊ Shift & pole introduction and swapping
- ◊ Rational Krylov
- ◊ Subspace iteration

Solution:

**SPOILER
ALERT**

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem:

- ◊ Small to medium-sized
- ◊ Regular
- ◊ All eigenvalues required
- ◊ No symmetry assumed

Tools:

- ◊ Hessenberg, Hessenberg form
- ◊ Shift & pole introduction and swapping
- ◊ Rational Krylov
- ◊ Subspace iteration

Solution:

- ◊ Method of QR-type driven by rational functions

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem:

- ◊ Small to medium-sized
- ◊ Regular
- ◊ All eigenvalues required
- ◊ No symmetry assumed

Tools:

- ◊ Hessenberg, Hessenberg form
- ◊ Shift & pole introduction and swapping
- ◊ Rational Krylov
- ◊ Subspace iteration

Solution:

- ◊ Method of QR-type driven by rational functions

Generalized eigenvalue problems

- ◊ Given A & B : $n \times n$ matrices, either \mathbb{R} or \mathbb{C}
- ◊ Computation of the triplets $(\alpha, \beta, \mathbf{x})$ that satisfy $\beta A\mathbf{x} = \alpha B\mathbf{x}$
- ◊ Procedure:
 1. Reduce the pencil to a *manageable* form
 2. Iterate to generalized Schur form
 3. Recover eigenvectors
- ◊ Make use of well-chosen unitary equivalences:
$$(\hat{A}, \hat{B}) = Q^*(A, B)Z \text{ preserves eigenvalues}$$

Generalized eigenvalue problems

- ◊ Given A & B : $n \times n$ matrices, either \mathbb{R} or \mathbb{C}
- ◊ Computation of the triplets $(\alpha, \beta, \mathbf{x})$ that satisfy $\beta A \mathbf{x} = \alpha B \mathbf{x}$
- ◊ Procedure:
 1. Reduce the pencil to a *manageable* form
 2. Iterate to generalized Schur form
 3. Recover eigenvectors
- ◊ Make use of well-chosen unitary equivalences:
$$(\hat{A}, \hat{B}) = Q^*(A, B)Z \text{ preserves eigenvalues}$$

Overview

In this talk we will discuss:

Numerical solution of the generalized eigenvalue problem $\beta A\mathbf{x} = \alpha B\mathbf{x}$:

- ◊ Small to medium-sized
- ◊ Regular
- ◊ All eigenvalues required
- ◊ No symmetry assumed

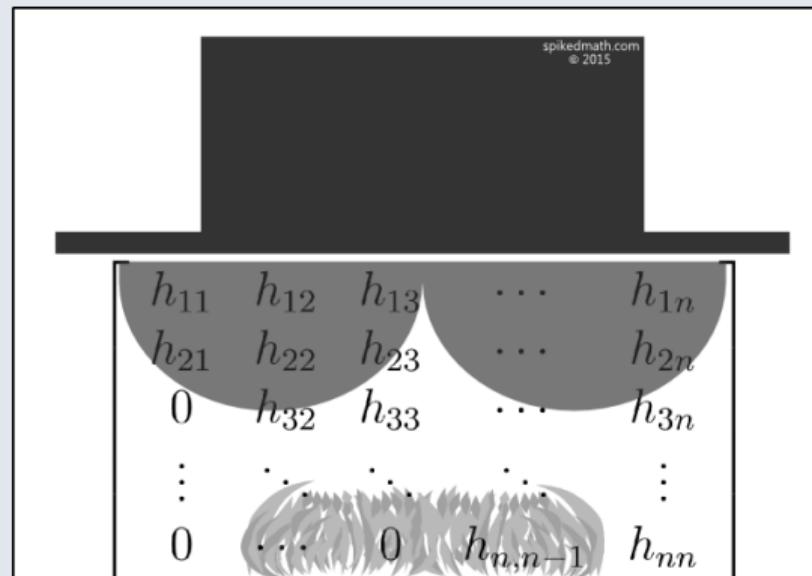
Tools:

- ◊ Hessenberg, Hessenberg form
- ◊ Shift & pole introduction and swapping
- ◊ Rational Krylov
- ◊ Subspace iteration

Solution:

- ◊ Method of QR-type driven by rational functions

Hessenberg, Hessenberg form



Upper Heisenberg Matrix

Hessenberg, Hessenberg form

A 7x10 grid of blue 'X' characters, arranged in seven rows and ten columns.

1

A 7x10 grid of blue 'X' characters, arranged in seven rows and ten columns.

A

B

Hessenberg, Hessenberg form

7

A

B

Hessenberg, Hessenberg form

$$\begin{matrix} \times & \times \\ \textcolor{green}{\times} & \times \\ \textcolor{green}{\times} & \times \\ \textcolor{green}{\times} & \times \\ \textcolor{green}{\times} & \times \\ \textcolor{green}{\times} & \times \\ \textcolor{green}{\times} & \times \\ \textcolor{green}{\times} & \times & \times & \end{matrix}$$

,

$$\begin{matrix} \times & \times \\ \textcolor{red}{\times} & \times \\ \textcolor{red}{\times} & \times \\ \textcolor{red}{\times} & \times \\ \textcolor{red}{\times} & \times \\ \textcolor{red}{\times} & \times \\ \textcolor{red}{\times} & \times \\ \textcolor{red}{\times} & \times & \end{matrix}$$

A

B

$$\text{poles } \Xi = \left(\frac{\textcolor{green}{\times}}{\textcolor{red}{\times}} \right) \subset \bar{\mathbb{C}}$$

Hessenberg, Hessenberg form

$\begin{matrix} \times & \times \\ \textcircled{1} & \times \\ \textcircled{2} & & \times \\ \textcircled{3} & & & \times & \times & \times & \times & \times & \times \\ \textcircled{4} & & & & \times & \times & \times & \times & \times \\ \textcircled{5} & & & & & \times & \times & \times & \times \\ \textcircled{6} & & & & & & \times & \times & \times \\ \textcircled{7} & & & & & & & \times & \end{matrix}$

,

$\begin{matrix} \times & \times \\ \textcircled{a} & & \times \\ \textcircled{b} & & & \times \\ \textcircled{c} & & & & \times & \times & \times & \times & \times & \times \\ \textcircled{d} & & & & & \times & \times & \times & \times & \times \\ \textcircled{e} & & & & & & \times & \times & \times & \times \\ \textcircled{f} & & & & & & & \times & \times & \times \\ \textcircled{g} & & & & & & & & \times & \end{matrix}$

$$\mathbf{A} \quad \mathbf{B}$$

poles $\Xi = \left(\frac{\textcircled{1}}{\textcircled{a}}, \frac{\textcircled{2}}{\textcircled{b}}, \dots \right) \subset \bar{\mathbb{C}}$

Introducing a shift

x	x	x	x	x	x	x	x
①	x	x	x	x	x	x	x
②	x	x	x	x	x	x	x
③	x	x	x	x	x	x	x
④	x	x	x	x			
⑤	x	x	x				
⑥	x	x					
⑦	x						

,

x	x	x	x	x	x	x	x
a	x	x	x	x	x	x	x
b	x	x	x	x	x	x	x
c	x	x	x	x	x	x	x
d	x	x	x	x			
e	x	x	x				
f	x	x					
g	x						

A

B

Introducing a shift

$\begin{bmatrix} \times & \times \\ \times & \times \end{bmatrix}$	\times	$\times \times \times \times \times \times \times$
	①	$\times \times \times \times \times \times \times$
	②	$\times \times \times \times \times \times$
	③	$\times \times \times \times \times$
	④	$\times \times \times \times$
	⑤	$\times \times \times$
	⑥	$\times \times$
	⑦	\times

A

$\begin{bmatrix} \times & \times \\ \times & \times \end{bmatrix}$	\times	$\times \times \times \times \times \times \times$
	a	$\times \times \times \times \times \times \times$
	b	$\times \times \times \times \times \times \times$
	c	$\times \times \times \times \times \times$
	d	$\times \times \times \times$
	e	$\times \times \times$
	f	$\times \times$
	g	\times

B

$$\boxed{\begin{array}{c|c} \times & \times \\ \textcircled{1} & \textcircled{a} \end{array} \neq \gamma !}$$

Introducing a shift



② $\times \times \times \times \times \times$
③ $\times \times \times \times \times$
④ $\times \times \times \times$
⑤ $\times \times \times$
⑥ $\times \times$
⑦ \times

,



b $\times \times \times \times \times \times$
c $\times \times \times \times \times$
d $\times \times \times \times$
e $\times \times \times$
f $\times \times$
g \times

A

B

Tools

Swapping poles

x	x	x	x	x	x	x	x
⊕	x	x	x	x	x	x	x
②	x	x	x	x	x	x	x
③	x	x	x	x	x	x	x
④	x	x	x	x	x	x	x
⑤	x	x	x	x	x	x	x
⑥	x	x	x	x	x	x	x
⑦	x	x	x	x	x	x	x

,

x	x	x	x	x	x	x	x
⊖	x	x	x	x	x	x	x
b	x	x	x	x	x	x	x
c	x	x	x	x	x	x	x
d	x	x	x	x	x	x	x
e	x	x	x	x	x	x	x
f	x	x	x	x	x	x	x
g	x	x	x	x	x	x	x

A

B

Swapping poles

Solve coupled Sylvester equation

(cfr. reordering Schur form [Kågström and Poromaa])

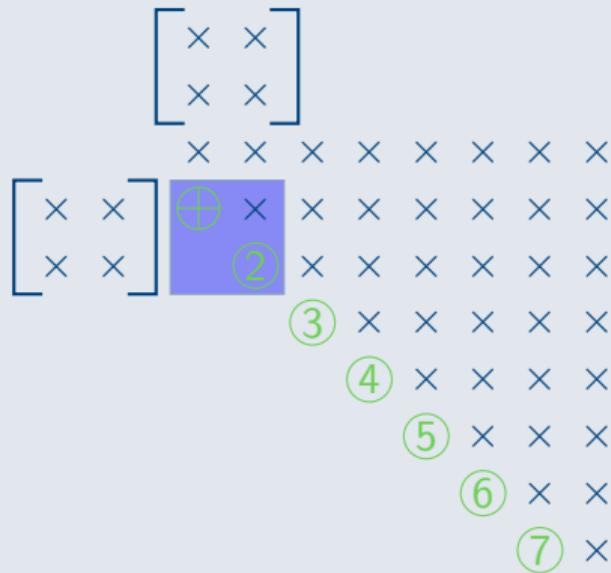
$$\frac{\oplus}{\ominus} \neq \frac{2}{b} !$$

$$\frac{\oplus}{\ominus} \neq \frac{0}{0} !$$

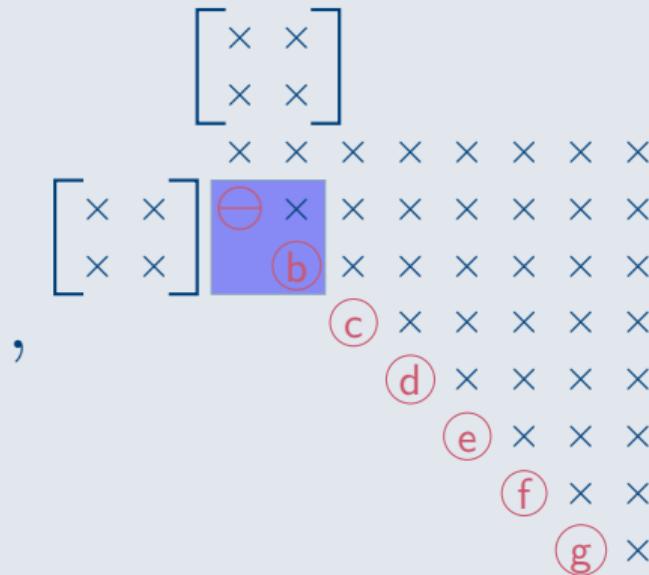
$$\frac{2}{b} \neq \frac{0}{0} !$$

$$\Rightarrow Q^* = \begin{bmatrix} \times & \times \\ \times & \times \end{bmatrix}, Z = \begin{bmatrix} \times & \times \\ \times & \times \end{bmatrix}$$

Swapping poles



A



B

Swapping poles

	⊗	⊗	×	×	×	×	×	×
②	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗
	⊕	⊗	⊗	⊗	⊗	⊗	⊗	⊗
③	×	×	×	×	×	×	×	×
④	×	×	×	×	×	×	×	×
⑤	×	×	×					
⑥	×							
⑦								

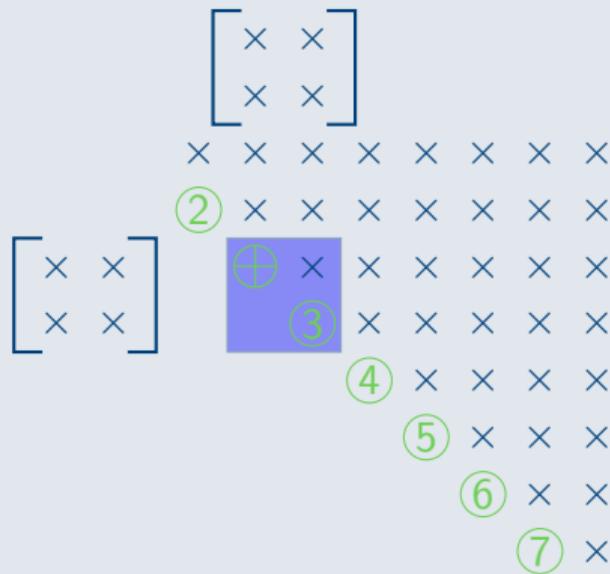
,

	⊗	⊗						
b	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗
	⊖	⊗	⊗	⊗	⊗	⊗	⊗	⊗
c								
d								
e								
f								
g								

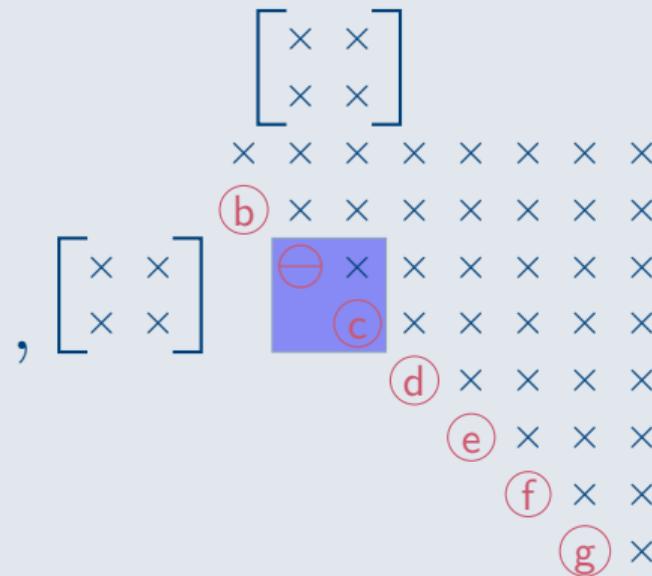
A

B

Swapping poles



A



B

Swapping poles

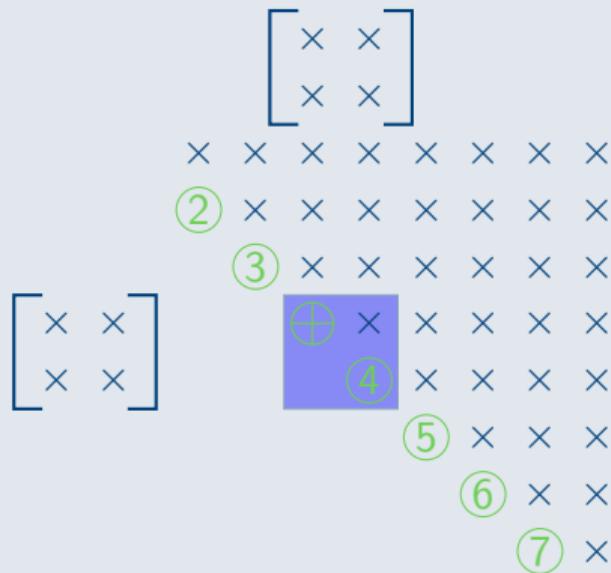
	\times	\otimes	\otimes	\times	\times	\times	\times	\times
②	\otimes	\otimes	\times	\times	\times	\times	\times	\times
③	\otimes							
	\oplus	\otimes						
④	\times	\times	\times	\times				
⑤	\times	\times	\times					
⑥	\times	\times						
⑦	\times							

,

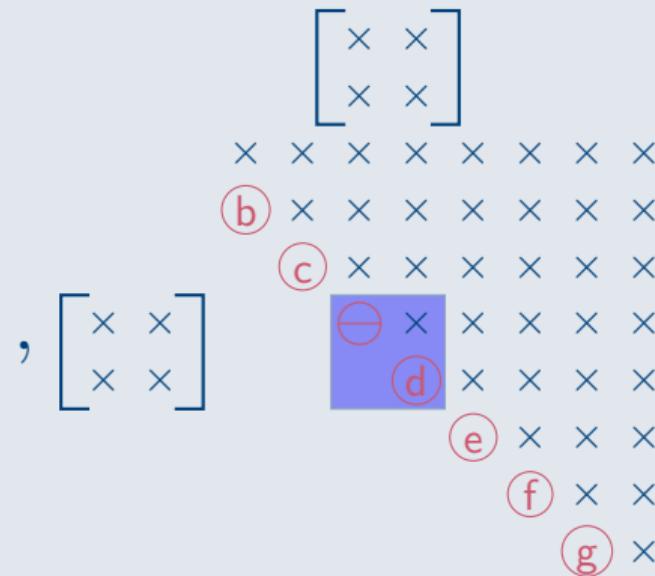
	\times	\otimes	\otimes	\times	\times	\times	\times	\times
b	\otimes	\otimes	\otimes	\times	\times	\times	\times	\times
c	\otimes							
	\ominus	\otimes						
d	\times	\times	\times	\times				
e	\times	\times	\times					
f	\times	\times						
g	\times							

A**B**

Swapping poles



A



B

Swapping poles

	x	x	⊗	⊗	x	x	x	x	x
②	x	⊗	⊗	x	x	x	x	x	x
③	⊗	⊗	x	x	x	x	x	x	x
④	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗
	⊕	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗
⑤	x	x	x						
⑥	x	x							
⑦	x								

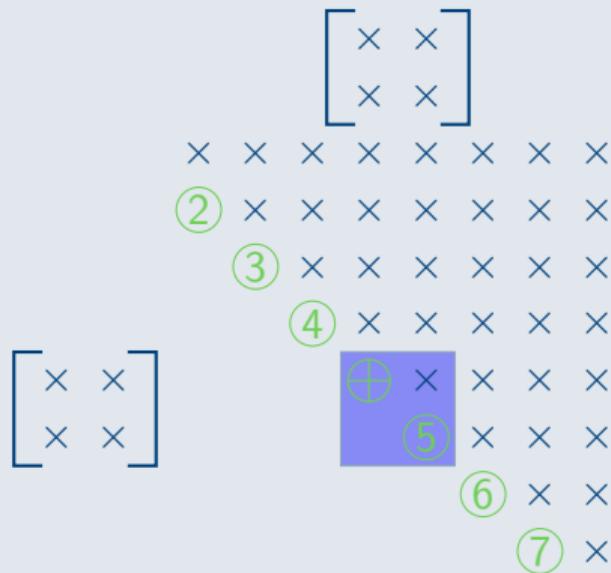
,

	x	x	⊗	⊗	x	x	x	x	x
b	x	⊗	⊗	x	x	x	x	x	x
c	⊗	⊗	x	x	x	x	x	x	x
d	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗
	⊖	⊗	⊗	⊗	⊗	⊗	⊗	⊗	⊗
e	x	x	x						
f	x	x							
g	x								

A

B

Swapping poles



A



B

Tools

Swapping poles

	\times	\times	\times	\otimes	\otimes	\times	\times	\times
②	\times	\times	\otimes	\otimes	\times	\times	\times	\times
③	\times	\otimes	\otimes	\times	\times	\times	\times	\times
④	\otimes	\otimes	\times	\times	\times	\times	\times	\times
⑤	\otimes							
	\oplus	\otimes						
⑥	\times	\times						
⑦	\times							

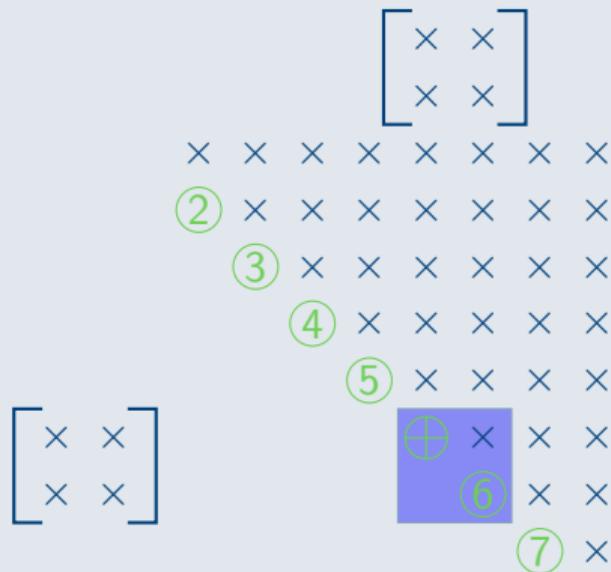
,

	\times	\times	\times	\otimes	\otimes	\times	\times	\times
b	\times	\times	\otimes	\otimes	\times	\times	\times	\times
c	\times	\otimes	\otimes	\times	\times	\times	\times	\times
d	\otimes	\otimes	\times	\times	\times	\times	\times	\times
e	\otimes							
	\ominus	\otimes						
f	\times	\times						
g	\times							

A

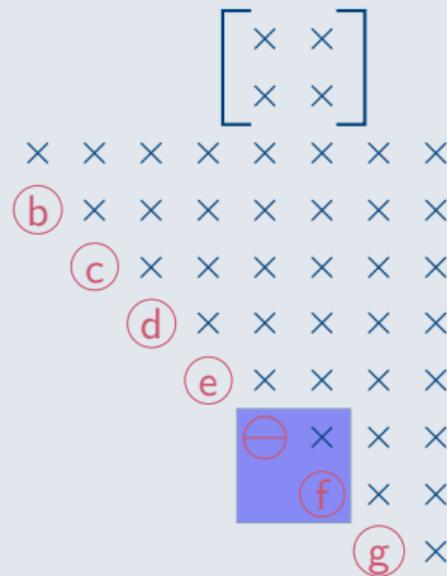
B

Swapping poles



A

,



B

Swapping poles

	\times	\times	\times	\times	\otimes	\otimes	\times	\times
②	\times	\times	\times	\otimes	\otimes	\times	\times	
③	\times	\times	\otimes	\otimes	\times	\times		
④	\times	\otimes	\otimes	\times	\times			
⑤	\otimes	\otimes	\times	\times				
⑥	\otimes	\otimes						
	\oplus	\otimes	\otimes					
⑦	\times							

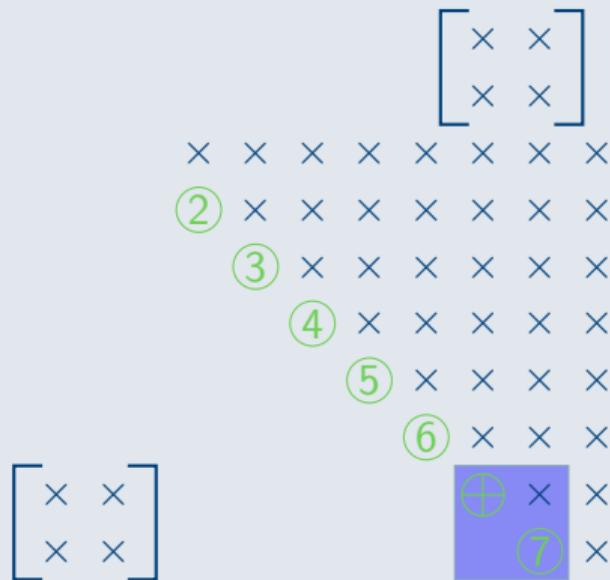
,

	\times	\times	\times	\times	\otimes	\otimes	\times	\times
b	\times	\times	\times	\otimes	\otimes	\times	\times	
c	\times	\times	\otimes	\otimes	\times	\times		
d	\times	\otimes	\otimes	\times	\times			
e	\otimes	\otimes	\times	\times				
f	\otimes	\otimes						
	\ominus	\otimes	\otimes					
g	\times							

A

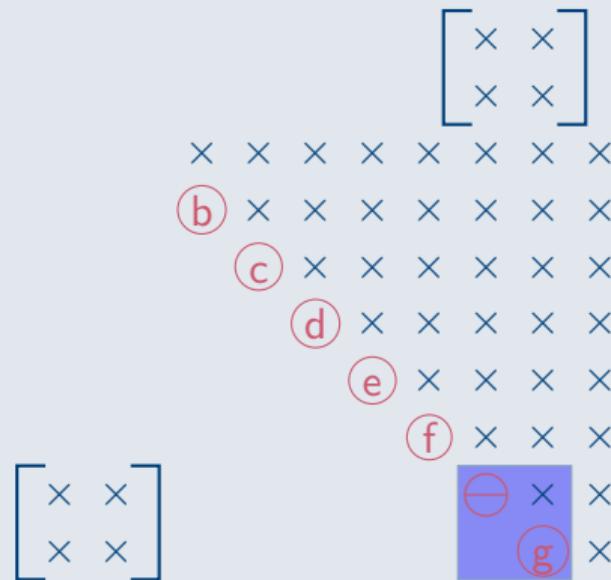
B

Swapping poles



A

,



B

Swapping poles

	\times	\times	\times	\times	\times	\otimes	\otimes	\times
②	\times	\times	\times	\times	\otimes	\otimes	\times	
③	\times	\times	\times	\otimes	\otimes	\times		
④	\times	\times	\otimes	\otimes	\times			
⑤	\times	\otimes	\otimes	\times				
⑥	\otimes	\otimes	\times					
⑦	\times	\otimes	\oplus	\otimes				

,

	\times	\times	\times	\times	\times	\otimes	\otimes	\times
b	\times	\times	\times	\times	\otimes	\otimes	\times	
c	\times	\times	\times	\otimes	\otimes	\times		
d	\times	\times	\otimes	\otimes	\times			
e	\times	\otimes	\otimes	\times				
f	\otimes	\otimes	\times					
g	\times	\otimes	\ominus	\otimes				

A

B

Introducing a pole

$$\begin{matrix} & \left[\begin{matrix} \times & \times \\ \times & \times \end{matrix} \right] \\ \times & \times \\ \textcircled{2} & \times \\ \textcircled{3} & \times \\ \textcircled{4} & \times \\ \textcircled{5} & \times \\ \textcircled{6} & \times \\ \textcircled{7} & \times & \times & \quad & \quad & \quad & \quad & \quad & \quad \\ & \oplus & \times & \quad & \quad & \quad & \quad & \quad & \quad \\ , & & & & & & & & \end{matrix} \quad \begin{matrix} & \left[\begin{matrix} \times & \times \\ \times & \times \end{matrix} \right] \\ \times & \times \\ \textcircled{b} & \times \\ \textcircled{c} & \times \\ \textcircled{d} & \times \\ \textcircled{e} & \times \\ \textcircled{f} & \times \\ \textcircled{g} & \times & \times & \quad \\ & \ominus & \times & \quad \end{matrix}$$

A

$$\boxed{\oplus \times \neq \gamma \ominus \times !}$$

B

Introducing a pole

	\times	\times	\times	\times	\times	\times	\otimes	\otimes
②	\times	\times	\times	\times	\times	\otimes	\otimes	
③	\times	\times	\times	\times	\otimes	\otimes		
④	\times	\times	\times	\otimes	\otimes			
⑤	\times	\times	\otimes	\otimes				
⑥	\times	\otimes	\otimes					
⑦	\otimes	\otimes						
⑧								

,

	\times	\times	\times	\times	\times	\times	\otimes	\otimes
b	\times	\times	\times	\times	\times	\otimes	\otimes	
c	\times	\times	\times	\times	\otimes	\otimes		
d	\times	\times	\times	\otimes	\otimes			
e	\times	\times	\otimes	\otimes				
f	\times	\otimes	\otimes					
g	\otimes	\otimes						
h								

A

B

The algorithm in a nutshell:

The algorithm in a nutshell:

1. Introduce shift at the top

The algorithm in a nutshell:

1. Introduce shift at the top
2. Swap it all the way down

The algorithm in a nutshell:

1. Introduce shift at the top
2. Swap it all the way down
3. Introduce pole at the end

The algorithm in a nutshell:

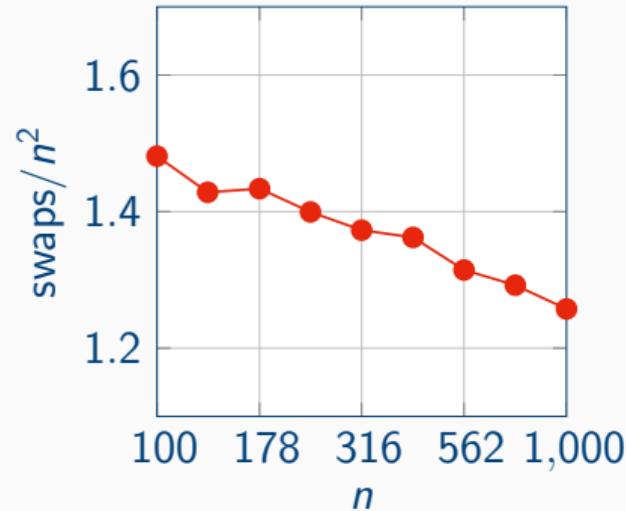
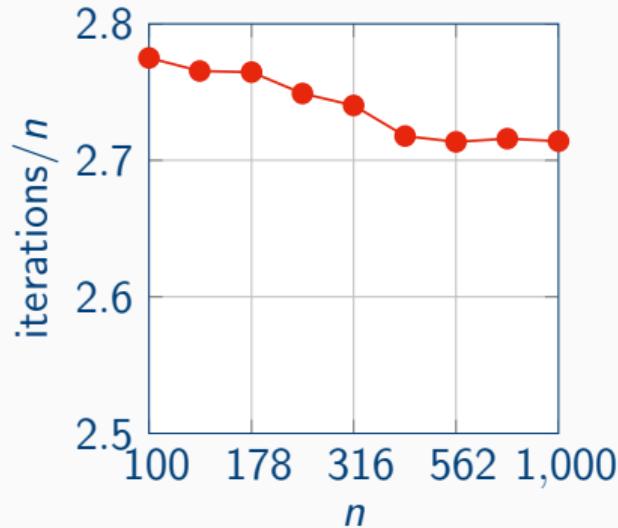
1. Introduce shift at the top
2. Swap it all the way down
3. Introduce pole at the end

Poles at ∞ ($\times = 0$) \rightarrow QZ method: Bulge exchange interpretation [Watkins]

Caution: shift $\notin \Xi$ to avoid slower convergence

Numerical example

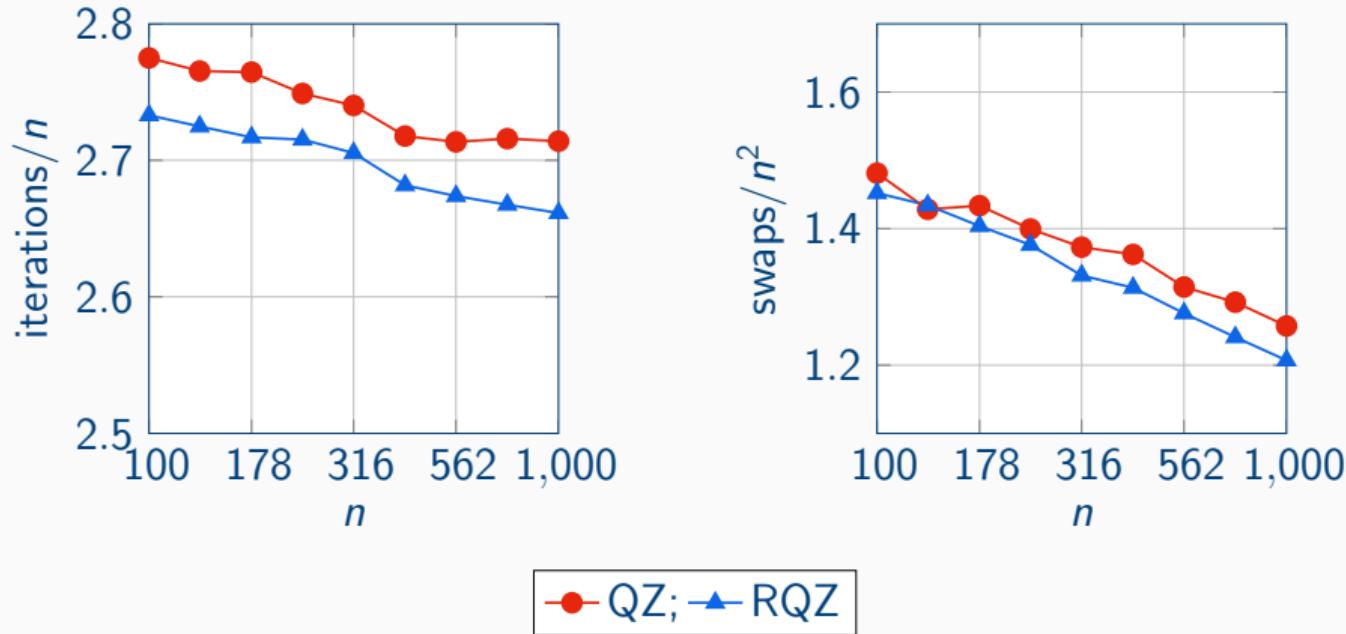
Is it worth it?



Data: 9 random matrix pairs, $n \in [100, 1000]$, reduced to H-T, averaged over 10 runs

Numerical example

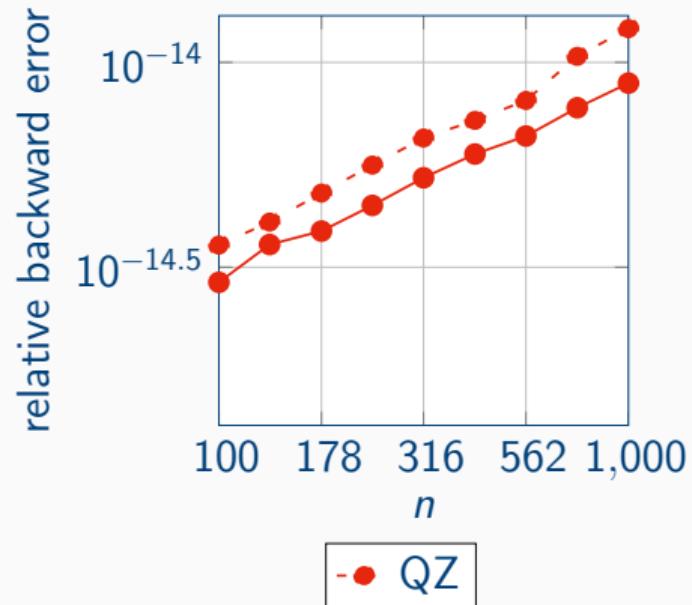
Is it worth it?



Data: 9 random matrix pairs, $n \in [100, 1000]$, reduced to H-T, averaged over 10 runs

Numerical example

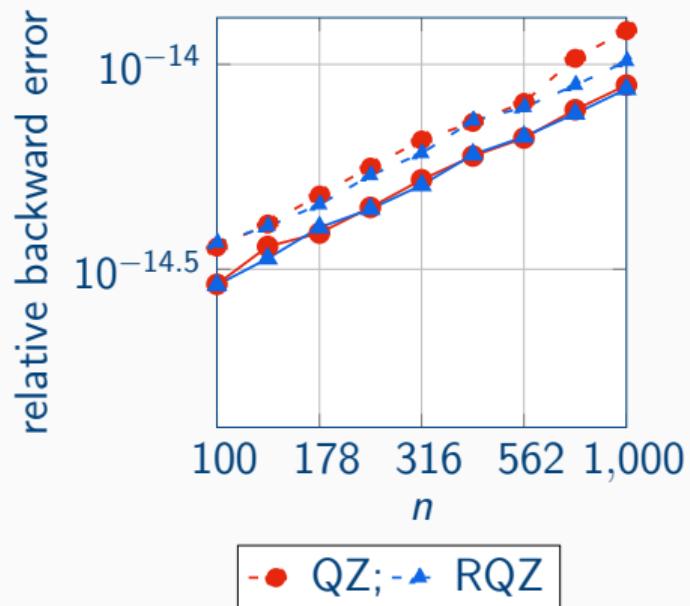
Is it worth it?



Data: 9 random matrix pairs, $n \in [100, 1000]$, reduced to H-T, averaged over 10 runs

Numerical example

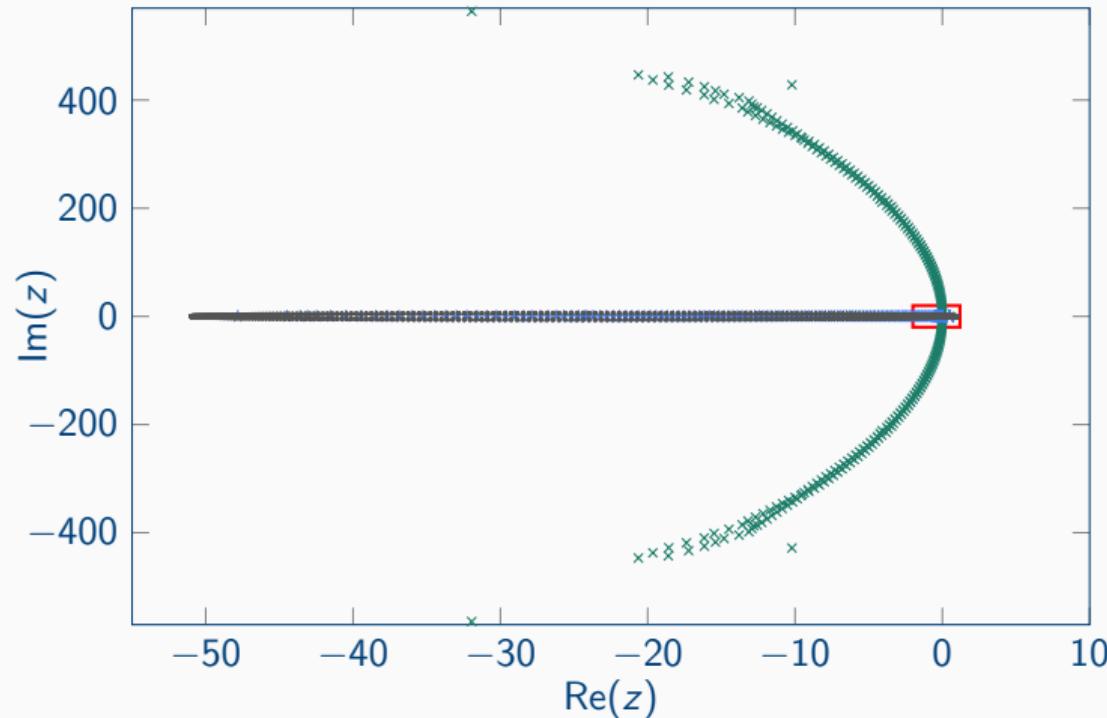
Is it worth it?



Data: 9 random matrix pairs, $n \in [100, 1000]$, reduced to H-T, averaged over 10 runs

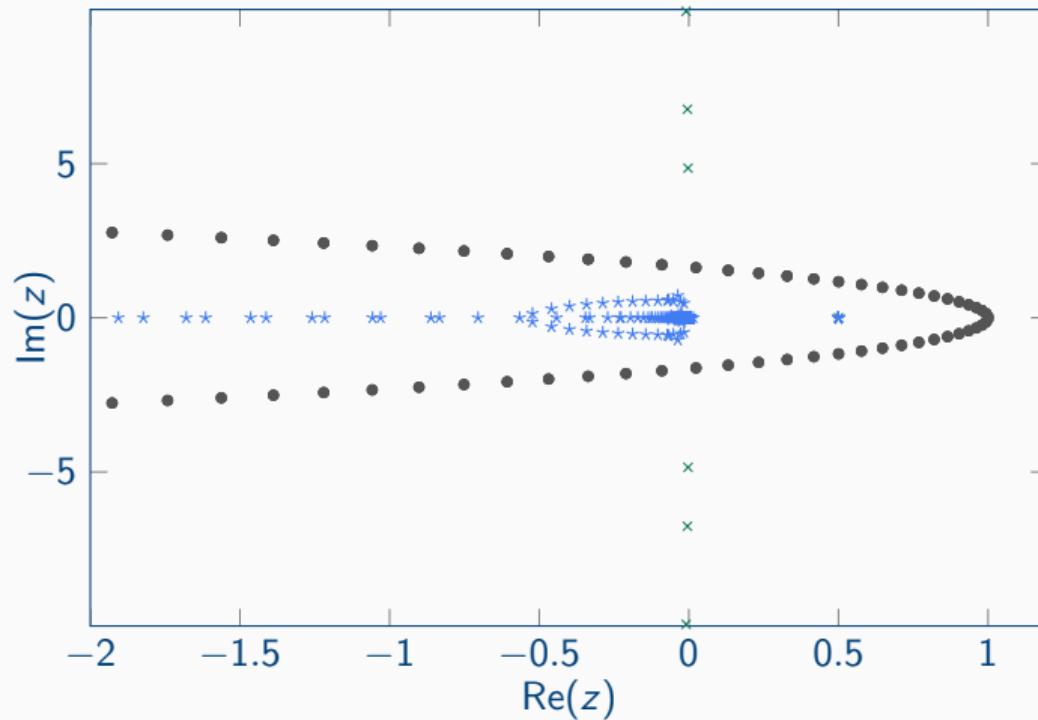
Numerical example 2: Reduction to Hessenberg, Hessenberg

Data: MHD matrix pair from MatrixMarket, $n = 1280$



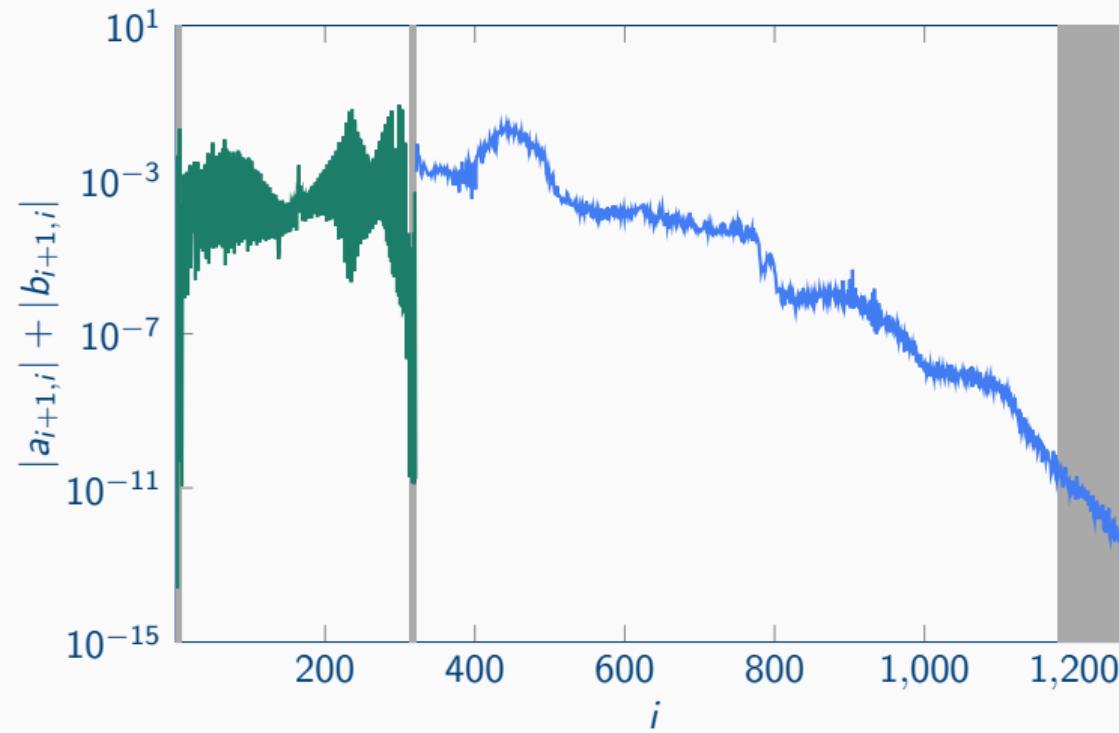
Numerical example 2: Reduction to Hessenberg, Hessenberg

Data: MHD matrix pair from MatrixMarket, $n = 1280$



Numerical example 2: Reduction to Hessenberg, Hessenberg

Data: MHD matrix pair from MatrixMarket, $n = 1280$



A word on the theory

Definition: Properness

The Hessenberg, Hessenberg pair (A, B) is called *proper* (or *irreducible*) if:

1.

$$\begin{array}{c} \times \\ \textcircled{1} \end{array} \neq \gamma \begin{array}{c} \times \\ \textcircled{a} \end{array}$$

2.

$$\begin{array}{c} \textcolor{green}{x} \\ \hline \textcolor{red}{x} \end{array} \neq \begin{array}{c} \textcolor{green}{0} \\ \hline 0 \end{array}$$

3.

$$\begin{array}{c} \oplus \\ \times \end{array} \neq \gamma \begin{array}{c} \ominus \\ \times \end{array}$$

A word on the theory

Why and how does RQZ work?

Krylov subspaces

1. Krylov subspace: $\mathcal{K}_i(M, \mathbf{v}) = \mathcal{R}(\mathbf{v}, M\mathbf{v}, \dots, M^{i-1}\mathbf{v})$

A word on the theory

Why and how does RQZ work?

Krylov subspaces

1. Krylov subspace: $\mathcal{K}_i(M, \mathbf{v}) = \mathcal{R}(\mathbf{v}, M\mathbf{v}, \dots, M^{i-1}\mathbf{v})$
2. rational Krylov subspace: $\mathcal{K}_i^{\text{rat}}(M, \mathbf{v}, \Xi = (\xi_1, \dots, \xi_{i-1})) = q_{\Xi}(M)^{-1} \mathcal{K}_i(M, \mathbf{v})$

A word on the theory

Why and how does RQZ work?

Krylov subspaces

1. Krylov subspace: $\mathcal{K}_i(M, \mathbf{v}) = \mathcal{R}(\mathbf{v}, M\mathbf{v}, \dots, M^{i-1}\mathbf{v})$
2. rational Krylov subspace: $\mathcal{K}_i^{\text{rat}}(M, \mathbf{v}, \Xi = (\xi_1, \dots, \xi_{i-1})) = q_{\Xi}(M)^{-1} \mathcal{K}_i(M, \mathbf{v})$

Theorem

If (A, B) is a proper Hessenberg pair with poles $(\xi_1, \dots, \xi_{n-1})$ then for $i = 1, \dots, n-1$:

$$\mathcal{K}_i^{\text{rat}}(AB^{-1}, \mathbf{e}_1, (\xi_1, \dots, \xi_{i-1})) = \mathcal{K}_i^{\text{rat}}(B^{-1}A, \mathbf{e}_1, (\xi_2, \dots, \xi_i)) = \mathcal{R}(\mathbf{e}_1, \dots, \mathbf{e}_i) = \mathcal{E}_i$$

A word on the theory

Why and how does RQZ work?

Theorem: Implicit Q (and Z)

Given a pair (A, B) , the matrices Q and Z that transform it to proper Hessenberg form,

$$(\hat{A}, \hat{B}) = Q^* (A, B) Z,$$

are determined *essentially unique* if $Q\mathbf{e}_1$ and the poles are fixed.

A word on the theory

Why and how does RQZ work?

Nested subspace iteration

An RQZ step with shift ϱ on a pencil with poles $(\xi_1, \dots, \xi_{n-1})$ and new pole ξ_n performs nested subspace iteration for $i = 1, \dots, n-1$ accelerated by

$$\mathcal{R}(\mathbf{q}_1, \dots, \mathbf{q}_i) = (A - \varrho B)(A - \xi_i B)^{-1} \mathcal{E}_i$$

$$\mathcal{R}(\mathbf{z}_1, \dots, \mathbf{z}_i) = (A - \xi_{i+1} B)^{-1}(A - \varrho B) \mathcal{E}_i.$$

A word on the theory

Why and how does RQZ work?

Nested subspace iteration

An RQZ step with shift ϱ on a pencil with poles $(\xi_1, \dots, \xi_{n-1})$ and new pole ξ_n performs nested subspace iteration for $i = 1, \dots, n-1$ accelerated by

$$\mathcal{R}(\mathbf{q}_1, \dots, \mathbf{q}_i) = (A - \varrho B)(A - \xi_i B)^{-1} \mathcal{E}_i$$

$$\mathcal{R}(\mathbf{z}_1, \dots, \mathbf{z}_i) = (A - \xi_{i+1} B)^{-1}(A - \varrho B) \mathcal{E}_i.$$

What does this mean?

- QR step with shift ϱ on entire space

A word on the theory

Why and how does RQZ work?

Nested subspace iteration

An RQZ step with shift ϱ on a pencil with poles $(\xi_1, \dots, \xi_{n-1})$ and new pole ξ_n performs nested subspace iteration for $i = 1, \dots, n-1$ accelerated by

$$\mathcal{R}(\mathbf{q}_1, \dots, \mathbf{q}_i) = (A - \varrho B)(A - \xi_i B)^{-1} \mathcal{E}_i$$

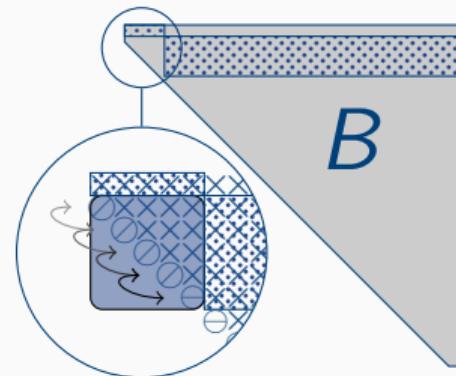
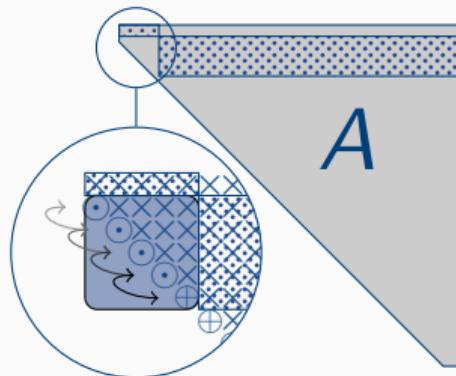
$$\mathcal{R}(\mathbf{z}_1, \dots, \mathbf{z}_i) = (A - \xi_{i+1} B)^{-1} (A - \varrho B) \mathcal{E}_i.$$

What does this mean?

- QR step with shift ϱ on entire space
- RQ steps with tightly packed shifts Ξ on selected subspaces

Further extensions

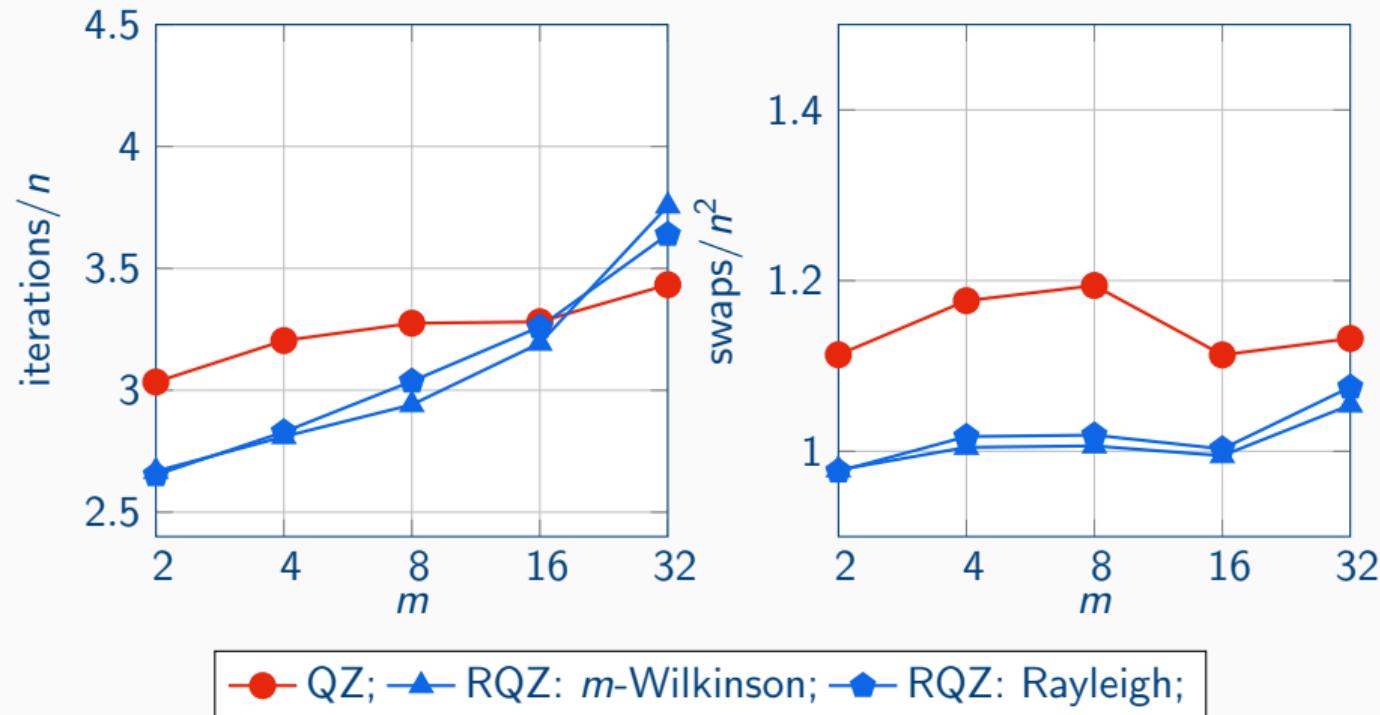
Tightly packed shifts



→ More cache efficient implementations (Level 3 BLAS)

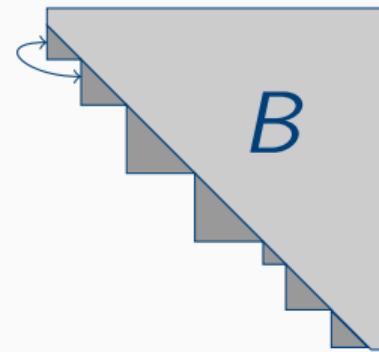
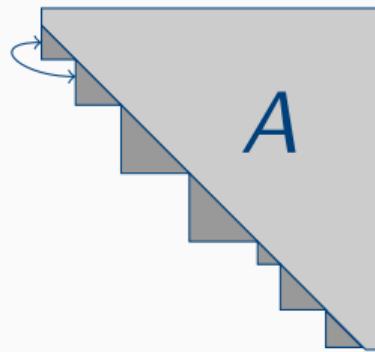
Further extensions

Tightly packed shifts



Further extensions

Block Hessenberg



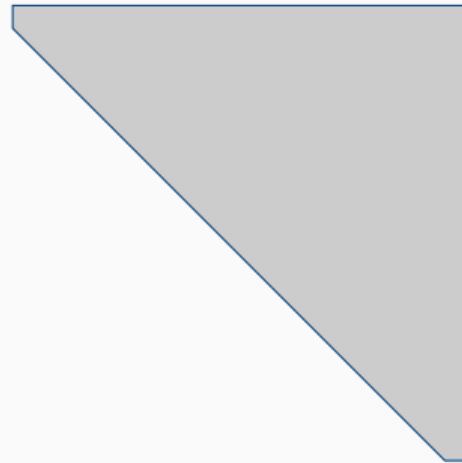
→ complex conjugate shifts and poles in real arithmetic for real pencils

Aggressive Early deflation

The performance of the QR algorithm can be significantly improved by an aggressive early deflation technique ([Braman, Byers and Mathias]) and similar techniques have been developed for the QZ method ([Kågström and Kressner]).

Aggressive Early deflation

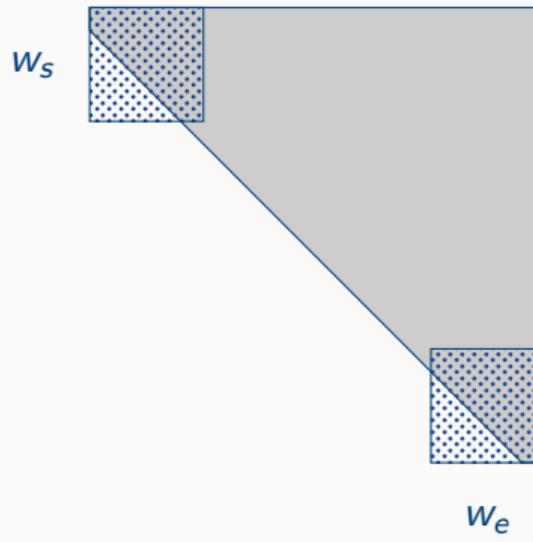
A and B



Further extensions

Aggressive Early deflation

A and B

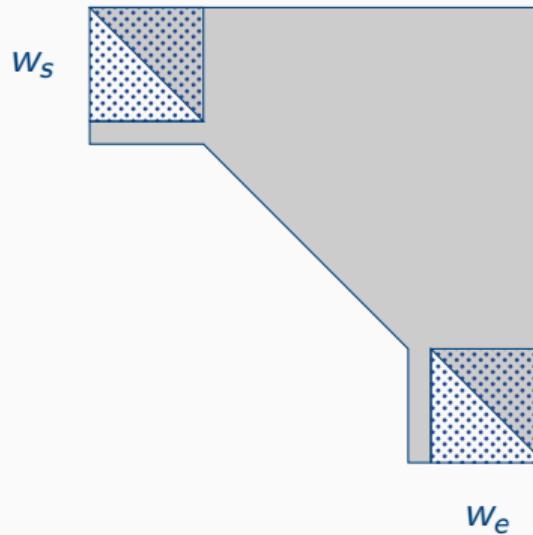


$$n \gg w_e > w_s$$

Further extensions

Aggressive Early deflation

A and B

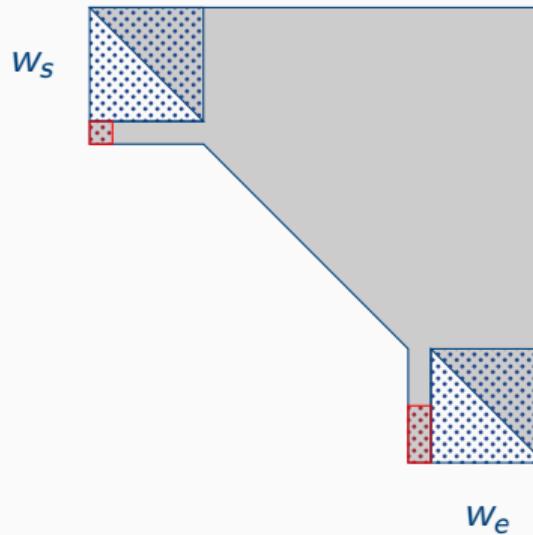


$$n \gg w_e > w_s$$

Further extensions

Aggressive Early deflation

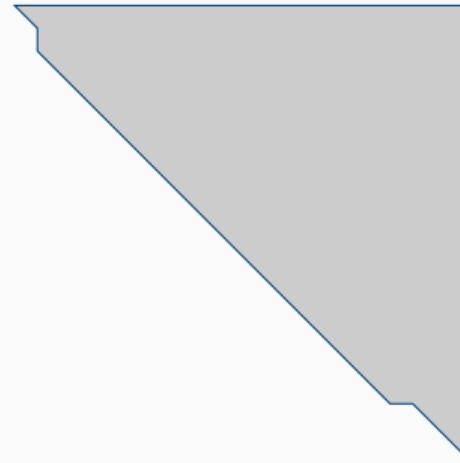
A and B



$$n \gg w_e > w_s$$

Aggressive Early deflation

A and B



Further extensions

Standard eigenvalue problems

- RQZ applies equivalence transformations on the pencil:

$$(\hat{A}, \hat{B}) = Q^*(A, B)Z$$

- Consequently we have two similarity transformations:

$$\hat{A}\hat{B}^{-1} = Q^*AB^{-1}Q \quad \text{and} \quad \hat{B}^{-1}\hat{A} = Z^*B^{-1}AZ$$

Further extensions

Standard eigenvalue problems

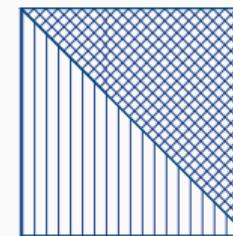
- RQZ applies equivalence transformations on the pencil:

$$(\hat{A}, \hat{B}) = Q^*(A, B)Z$$

- Consequently we have two similarity transformations:

$$\hat{A}\hat{B}^{-1} = Q^*AB^{-1}Q \quad \text{and} \quad \hat{B}^{-1}\hat{A} = Z^*B^{-1}AZ$$

$$\hat{A}\hat{B}^{-1} - \text{diag}(\alpha, \xi_1, \dots, \xi_{n-1}) =$$



Further extensions

Standard eigenvalue problems

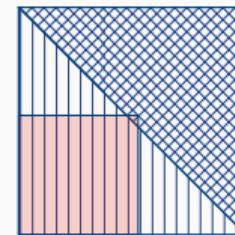
- RQZ applies equivalence transformations on the pencil:

$$(\hat{A}, \hat{B}) = Q^*(A, B)Z$$

- Consequently we have two similarity transformations:

$$\hat{A}\hat{B}^{-1} = Q^*AB^{-1}Q \quad \text{and} \quad \hat{B}^{-1}\hat{A} = Z^*B^{-1}AZ$$

$$\hat{A}\hat{B}^{-1} - \text{diag}(\alpha, \xi_1, \dots, \xi_{n-1}) =$$



Further extensions

Standard eigenvalue problems

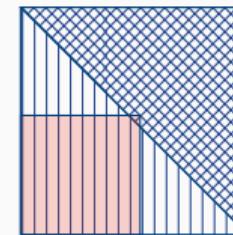
- RQZ applies equivalence transformations on the pencil:

$$(\hat{A}, \hat{B}) = Q^*(A, B)Z$$

- Consequently we have two similarity transformations:

$$\hat{A}\hat{B}^{-1} = Q^*AB^{-1}Q \quad \text{and} \quad \hat{B}^{-1}\hat{A} = Z^*B^{-1}AZ$$

$$\hat{A}\hat{B}^{-1} - \text{diag}(\alpha, \xi_1, \dots, \xi_{n-1}) =$$



Extended Hessenberg + diagonal = rational Hessenberg

→ my connection with this mini-symposium

Rational Krylov method

The connection between RQZ and the rational Krylov method can be used:

- to compute the Ritz values from the Hessenberg, Hessenberg recurrence pencil
- to filter and restart the rational Krylov method

Rational Krylov method

$$A V_{k+1} \underline{G}_k = B V_{k+1} \underline{H}_k$$

with:

- $\mathcal{R}(V_{k+1}) = \mathcal{K}_{k+1}^{\text{rat}}(AB^{-1}, \mathbf{v}, \Xi_{1:k})$
- $(\underline{H}_k, \underline{G}_k)$ the Hessenberg, Hessenberg recurrence pencil

Applying an RQZ step with shift ϱ , we get $\mathcal{K}_k^{\text{rat}}(AB^{-1}, \hat{\mathbf{v}}, \Xi_{2:k})$ with:

$$\hat{\mathbf{v}} = (A - \xi_1 B)^{-1} (A - \varrho B) \mathbf{v}$$

Conclusions

1. RQZ is a generalization of QZ
2. Implicit rational subspace iteration is promising
3. New shift and pole strategies can be a powerful tool to compute invariant subspaces

Further reading:

arXiv:1802.04094

<http://numa.cs.kuleuven.be/software/rqz/>

Conclusions

1. RQZ is a generalization of QZ
2. Implicit rational subspace iteration is promising
3. New shift and pole strategies can be a powerful tool to compute invariant subspaces

Further reading:

arXiv:1802.04094

<http://numa.cs.kuleuven.be/software/rqz/>

Thank you for your attention!